# Optimal Rotation Angles for Quasi-Orthogonal Space-Time Codes with PSK Modulation 

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#### Abstract

Full diversity of quasi-orthogonal space-time block codes with complex signals could be achieved by constellation rotation. In this letter, we derive the optimal rotation angle in the sense of maximizing coding gain for quasi-orthogonal codes with phase-shift keying (PSK) modulation using a geometrybased approach. We also prove that the rotated PSK signals with an even constellation size have higher coding gains (or diversity products) than those with an odd constellation size.


Index Terms-Quasi-orthogonal space-time codes, constellation rotation, phase-shift keying (PSK).

## I. Introduction

THE transmit diversity scheme proposed in [1] is a simple and effective orthogonal space-time block code of rate one for systems with two transmit antennas. The orthogonal design was then extended to systems with an arbitrary number of transmit antennas [2]. Rate-one real orthogonal codes are available from [2], but complex orthogonal design with transmission rate one does not exist for more than two transmit antennas (see [2, Theorem 5.4.2]). Quasi-orthogonal spacetime block codes that provide rate one but partial diversity were studied in [3], [4] for four transmit antennas, and in [5] for higher number of transmit antennas. To achieve full diversity for quasi-orthogonal codes, constellation rotation schemes were proposed in [6], [7]. Optimal rotation angles in the sense of maximizing coding gain for quadrature amplitude modulation were addressed in [7]. In [8], optimal rotation angles for PSK with an even constellation size were derived.

In this letter, we derive, through a geometry-based approach, the optimal rotation angles for quasi-orthogonal codes with any PSK modulation. The independent work [9] also addressed the optimal rotation angles using a completely different approach. In addition to the optimal rotation angles, we also prove that coding gain for even-sized constellations is higher than that for odd-sized constellations, which was observed but not proved in [9].

## II. Optimal constellation rotation for PSK

We focus on the scheme given in [4] for systems with four transmit antennas for which the code matrix is expressed as $\boldsymbol{C}=\left[\begin{array}{cccc}x_{1} & x_{2} & x_{3} & x_{4} \\ -x_{2}^{*} & x_{1}^{*} & -x_{4}^{*} & x_{3}^{*} \\ x_{3} & x_{4} & x_{1} & x_{2} \\ -x_{4}^{*} & x_{3}^{*} & -x_{2}^{*} & x_{1}^{*}\end{array}\right]$, and $\boldsymbol{C}^{H} \boldsymbol{C}=\left[\begin{array}{cccc}a & 0 & b & 0 \\ 0 & a & 0 & b \\ b & 0 & a & 0 \\ 0 & b & 0 & a\end{array}\right]$,

Manuscript received January 8, 2005. The associate editor coordinating the review of this letter and approving it for publication was Prof. George K. Karagiannidis.

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Digital Object Identifier 10.1109/LCOMM.2005.08021.
where $(\cdot)^{*}$ denotes complex conjugate, $(\cdot)^{H}$ denotes conjugate transpose, $a=\sum_{i=1}^{4}\left|x_{i}\right|^{2}$, and $b=x_{1} x_{3}^{*}+x_{3} x_{1}^{*}+x_{2} x_{4}^{*}+$ $x_{4} x_{2}^{*}$. Analysis for other quasi-orthogonal codes (e.g., the code given in [3]) is similar.

The maximum likelihood (ML) decision metric for this code can be written as the sum of two independent terms $f_{1}\left(x_{1}, x_{3}\right)+f_{2}\left(x_{2}, x_{4}\right)$. Thus, the minimization for ML decoding can be done separately on these two terms. For the code example chosen above, let us consider $\left(x_{1}, x_{3}\right)$ and let $x_{2}=x_{4}=0$ [6] in calculating the optimal rotation angles.

Let $\mathcal{A}$ be a PSK constellation of size $Q$, where $Q$ could be even $(Q=2 n, n>0$ is an integer) or odd $(Q=$ $2 n-1, n>1$ ), and $\mathcal{B}$ be the rotated constellation of $\mathcal{A}$ expressed as $\mathcal{B}=e^{j \phi} \mathcal{A}$, where $\phi$ represents the rotation angle. Also let $x_{1}, \tilde{x}_{1} \in \mathcal{A}$ and $x_{3}, \tilde{x}_{3} \in \mathcal{B}$, where $\left(x_{1}, x_{3}\right) \neq$ ( $\tilde{x}_{1}, \tilde{x}_{3}$ ). Maximizing coding gain is equivalent to maximizing $\left|\operatorname{det}\left[\Delta_{C\left(x_{1}, 0, x_{3}, 0\right)}^{H} \Delta_{C\left(x_{1}, 0, x_{3}, 0\right)}\right]\right|$, where $|\cdot|$ denotes the absolute value and $\Delta_{C\left(x_{1}, 0, x_{3}, 0\right)}$ is the difference matrix given as $\Delta_{C\left(x_{1}, 0, x_{3}, 0\right)}=\boldsymbol{C}_{\left(x_{1}, 0, x_{3}, 0\right)}-\boldsymbol{C}_{\left(\tilde{x}_{1}, 0, \tilde{x}_{3}, 0\right)}$. This is also equivalent to maximizing the minimum $\zeta$-distance between constellations $\mathcal{A}$ and $\mathcal{B}$ expressed as [6, Eq. (23)]

$$
\begin{equation*}
d_{\min , \zeta}(\mathcal{A}, \mathcal{B}) \triangleq \underbrace{\min }_{\left(x_{1}, x_{3}\right) \neq\left(\tilde{x}_{1}, \tilde{x}_{3}\right)}\left|\left(x_{1}-\tilde{x}_{1}\right)^{2}-\left(x_{3}-\tilde{x}_{3}\right)^{2}\right|^{\frac{1}{2}} \tag{1}
\end{equation*}
$$

The set of values of $\left(x_{1}-\tilde{x}_{1}\right)^{2}$ with $x_{1}, \tilde{x}_{1} \in \mathcal{A}$ and $\left(x_{3}-\tilde{x}_{3}\right)^{2}$ with $x_{3}, \tilde{x}_{3} \in \mathcal{B}$ constitute, respectively, constellation $\mathcal{X}$ and constellation $\mathcal{Y}=e^{j 2 \phi} \mathcal{X}$. Note that both $\mathcal{X}$ and $\mathcal{Y}$ include the origin. The optimal rotation angle in the sense of maximizing coding gain must maximize the minimum distance between any point from $\mathcal{X}$ and any point from $\mathcal{Y}$, except the origin.

Before calculating the optimal rotation angles, let us examine the properties of $\mathcal{X}$. Examples of an even-sized constellation $\mathcal{A}$ (16PSK) and an odd-sized constellation $\mathcal{A}$ (7PSK) are shown in Fig. 1 and Fig. 2, respectively.

The corresponding new constellations $\mathcal{X}$ formed by the set of values of $\left(x_{1}-\tilde{x}_{1}\right)^{2}$ with $x_{1}, \tilde{x}_{1} \in \mathcal{A}$ are shown in Fig. 3 and Fig. 4, respectively. In the following discussion, we exclude the origin in $\mathcal{X}$ because if $x_{1}-\tilde{x}_{1}=0$ or $x_{3}-\tilde{x}_{3}=$ $0,\left|\left(x_{1}-\tilde{x}_{1}\right)^{2}-\left(x_{3}-\tilde{x}_{3}\right)^{2}\right|^{1 / 2}$ becomes a constant for any rotation angle. Let $P=Q / 2$ when $Q$ is an even number and $P=(Q-1) / 2$ when $Q(Q \geq 3)$ is an odd number. In the following proofs, we focus on the case when $Q$ is an even number. The proof when $Q$ is an odd number is similar.
Properties of $\mathcal{X}$ :

1) All points in $\mathcal{X}$ can be divided into $P$ groups, $S_{1}, S_{2}, \cdots, S_{P}$, based on their relative magnitudes in ascending order as $\left|S_{1}\right|<\left|S_{2}\right|<\cdots<\left|S_{P}\right|$.


Fig. 1. An example of an even-sized constellation $\mathcal{A}$ : 16PSK.


Fig. 2. An example of an odd-sized constellation $\mathcal{A}$ : 7PSK.

Proof: It is easy to see from Fig. 1 that for all possible combinations of $x_{1}, \tilde{x}_{1} \in \mathcal{A}$ and $x_{1} \neq \tilde{x}_{1},\left|x_{1}-\tilde{x}_{1}\right|$ has $Q / 2$ distinct values. Additionally, the group with the minimum magnitude is formed by two adjacent points in $\mathcal{A}$ such as $x_{1}=p_{1}$ and $\tilde{x}_{1}=p_{2}$. The minimum magnitude is thus obtained to be $\left|S_{1}\right|=\left|p_{2}-p_{1}\right|^{2}=$ $2-2 \cos (2 \pi / Q)$. Obviously, we have $d_{\min , \zeta}(\mathcal{A}, \mathcal{B}) \leq$ $\left|S_{1}\right|^{1 / 2}$.
2) The distance between any two points from two different groups is greater than or equal to $\left|S_{1}\right|$ when $Q$ is an even number and greater than $\left|S_{1}\right|$ when $Q$ is an odd number.
Proof: Consider the worst case where two points from two adjacent groups have the same phase. The distance between these two points is $\left|S_{i}\right|-\left|S_{i-1}\right|$. From Fig. 1 , it is easy to see that sides $\sqrt{\left|S_{1}\right|}, \sqrt{\left|S_{i-1}\right|}$, and $\sqrt{\left|S_{i}\right|}$ always constitute a triangle. Additionally, the angle opposite $\sqrt{\left|S_{i}\right|}$ is an obtuse angle or right angle (e.g., $\angle p_{1} p_{7} p_{8}$ is an obtuse angle and $\angle p_{1} p_{8} p_{9}$ is a right angle). Thus, we have $\left|S_{i}\right|-\left|S_{i-1}\right| \geq\left|S_{1}\right|$ with equality if and only if the triangle is a right triangle.


Fig. 3. Constellation $\mathcal{X}$ formed by $\left(x_{1}-\tilde{x}_{1}\right)^{2}$, where $x_{1}, \tilde{x}_{1} \in \mathcal{A}$ as shown in Fig. 1.


Fig. 4. Constellation $\mathcal{X}$ formed by $\left(x_{1}-\tilde{x}_{1}\right)^{2}$, where $x_{1}, \tilde{x}_{1} \in \mathcal{A}$ as shown in Fig. 2.

For the example shown in Fig. 3, the distance between the two outer circles equals the radius of the inner circle.
3) There are exactly $Q / 2(Q=2 n)$ or $Q(Q=2 n-1)$ different points in each of the $P$ groups, $S_{1}, S_{2}, \cdots, S_{P}$, and the phases of all points in any group are uniformly distributed between 0 and $2 \pi$.
Proof: We prove using group $S_{1}$ as an example. The same method applies to other groups. All points in $S_{1}$ are listed as

$$
\begin{aligned}
\left(p_{2}-p_{1}\right)^{2} & =e^{j \pi}\left(e^{-j 2 \pi / Q}-1\right)^{2} \\
& \vdots \\
\left(p_{Q}-p_{Q-1}\right)^{2} & =e^{-j 4 \pi(Q-2) / Q} \cdot e^{j \pi}\left(e^{-j 2 \pi / Q}-1\right)^{2} \\
\left(p_{1}-p_{Q}\right)^{2} & =e^{-j 4 \pi(Q-1) / Q} \cdot e^{j \pi}\left(e^{-j 2 \pi / Q}-1\right)^{2}
\end{aligned}
$$

It is found that the phase difference between any two adjacent points in $S_{1}$ is $4 \pi / Q$ when $Q$ is an even number, and is $2 \pi / Q$ when $Q$ is an odd number.

## A. The Optimal Rotation Angles

According to properties 1-3 of constellation $\mathcal{X}$, the optimal rotation angle for quasi-orthogonal codes with PSK modulation is $\pi / Q$ when $Q$ is even and $\pi /(2 Q)$ when $Q$ is odd.

Proof: If we let $\phi=\pi / Q$, constellation $\mathcal{Y}$ is related to constellation $\mathcal{X}$ as $\mathcal{Y}=e^{j 2 \phi} \mathcal{X}=e^{j 2 \pi / Q} \mathcal{X}$. According to property 3 of $\mathcal{X}$, the minimum distance between any two points $s \in S_{i}$ and $\tilde{s} \in e^{j 2 \pi / Q} S_{i}$ is maximized. Additionally, from property 2 of $\mathcal{X}$, the minimum distance between $\alpha \in S_{i}$ and $\beta \in e^{j 2 \pi / Q} S_{j}, i \neq j$, is always greater than or equal to $\left|S_{1}\right|$. Therefore, the minimum $\zeta$-distance is determined by $\left|S_{1}\right|$, and $\phi=\pi / Q$ is the optimal rotation angle that maximizes $d_{\min , \zeta}(\mathcal{A}, \mathcal{B})$ when $Q$ is even. Following a similar procedure, we can prove the conclusion for odd-sized constellations. Note that rotation with $\phi=\pi / Q$ for oddsized constellations does not even provide full diversity ( [10], Theorem 2.2). This is clear seen from Fig. 4: if $\mathcal{Y}=e^{j 2 \pi / Q} \mathcal{X}$, then $\mathcal{Y}$ and $\mathcal{X}$ overlap and $\left|\left(x_{1}-\tilde{x}_{1}\right)^{2}-\left(x_{3}-\tilde{x}_{3}\right)^{2}\right|^{\frac{1}{2}}$ could be zero.

The corresponding optimal minimum $\zeta$-distance is

$$
\begin{aligned}
& d_{\min , \zeta}(\mathcal{A}, \mathcal{B})= \\
& \begin{cases}\min \left(\left.\left.|2| S_{1}\right|^{2}\left(1-\cos \left(\frac{2 \pi}{Q}\right)\right)\right|^{\frac{1}{4}},\left|S_{1}\right|^{\frac{1}{2}}\right), & Q=2 n \\
\min \left(\left.\left.|2| S_{1}\right|^{2}\left(1-\cos \left(\frac{\pi}{Q}\right)\right)\right|^{\frac{1}{4}},\left|S_{1}\right|^{\frac{1}{2}}\right), & Q=2 n-1\end{cases}
\end{aligned}
$$

The above expression can be simplified as

$$
d_{\min , \zeta}(\mathcal{A}, \mathcal{B})= \begin{cases}2 \sin \left(\frac{\pi}{Q}\right), & Q=2,4  \tag{2}\\ \sqrt{8 \sin ^{3}\left(\frac{\pi}{Q}\right),} & Q=2 n \geq 6 \\ \sqrt{8 \sin \left(\frac{\pi}{2 Q}\right) \sin ^{2}\left(\frac{\pi}{Q}\right),} & Q=2 n-1\end{cases}
$$

For the specific case of 8PSK ( $Q=8$ ), the optimal rotation angle based on the conclusion in this letter is $\phi=\pi / 8$ and the minimum $\zeta$-distance is 0.6696 , which are the same as the results obtained via computer search in [7].

## B. Relative Coding Gain between $Q=2 n$ and $Q=2 n-1$

Proposition 1: The cases of $Q=2 n, n \geq 3$, have larger optimal minimum $\zeta$-distances than those of $Q=2 n-1$. Mathematically, Proposition 1 is expressed as

$$
\begin{equation*}
\sin ^{3}\left(\frac{\pi}{2 n}\right)-\sin \left(\frac{\pi}{2(2 n-1)}\right) \sin ^{2}\left(\frac{\pi}{2 n-1}\right)>0 \tag{3}
\end{equation*}
$$

Before proving Proposition 1, let us prove the following inequality:

$$
\begin{equation*}
x>\sin (x)>x-x^{3} / 6, \text { for } x>0 \tag{4}
\end{equation*}
$$

Let $f_{1}(x)=x-\sin (x)$ and $f_{2}(x)=\sin (x)-x+x^{3} / 6$. Obviously, $\left.f_{1}(x)\right|_{x=0}=0$ and $f_{1}^{\prime}(x) \geq 0, \forall x>0$, where $(\cdot)^{\prime}$
denotes derivative. Thus, the left inequality of (4) follows. The derivative of $f_{2}(x)$ with respect to $x$ is written as

$$
\begin{align*}
f_{2}^{\prime}(x) & =\cos (x)-1+\frac{x^{2}}{2} \\
& =2\left[\left(\frac{x}{2}\right)^{2}-\sin ^{2}\left(\frac{x}{2}\right)\right]>0 \tag{5}
\end{align*}
$$

Thus, the right inequality of (4) follows.
Let $\xi_{1}=\left[\frac{\pi}{2 n}-\frac{(\pi /(2 n))^{3}}{6}\right]^{3}$ and $\xi_{2}=\frac{\pi}{2(2 n-1)}\left(\frac{\pi}{2 n-1}\right)^{2}$. For $n \geq 3$

$$
\xi_{1}-\xi_{2}=(n+0.92)(n-0.36)(n-2.98)>0
$$

Notice that from (4) we have

$$
\begin{aligned}
& \xi_{1}<\sin ^{3}\left(\frac{\pi}{2 n}\right) \\
& \xi_{2}>\sin \left(\frac{\pi}{2(2 n-1)}\right) \sin ^{2}\left(\frac{\pi}{2 n-1}\right)
\end{aligned}
$$

Thus, Proposition 1 expressed in (3) is proved.

## III. CONCLUSION

We have derived, using a geometry-based method, the optimal constellation rotation angles for quasi-orthogonal spacetime block codes for four-antenna systems with PSK modulation. Through constellation rotation with these rotation angles, the coding gain is maximized and full diversity of quasi-orthogonal codes is achieved. We have also proved that PSK signals with a constellation size $Q=2 n$ have larger optimal minimum $\zeta$-distances than those with a constellation size $Q=2 n-1(n \geq 3)$.

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