Architecture: Caching Issues in Performance

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Problem: The Path Between a CPU Chip and Off-chip Memory is Slow

This path is relatively slow, forcing the CPU to wait for up to 200 clock cycles just to do a store to, or a load from, memory. Depending on your CPU's ability to process instructions out-of-order, it might go idle during this time. This is a huge performance hit!

Solution: Hierarchical Memory Systems, or “Cache”

The solution is to add intermediate memory systems. The one closest to the CPU is small and fast. The memory systems get slower and larger as they get farther away from the CPU.

In computer science, a cache is a collection of data duplicating original values stored elsewhere or computed earlier, where the original data is expensive to fetch (due to longer access time) or to compute, compared to the cost of reading the cache. In other words, a cache is a temporary storage area where frequently accessed data can be stored for rapid access. Once the data is stored in the cache, future use can be made by accessing the cached copy rather than re-fetching or recomputing the original data, so that the average access time is shorter. Cache, therefore, helps expedite data access that the CPU would otherwise need to fetch from main memory.

-- Wikipedia

Cache and Memory are Named by “Distance Level” from the CPU

L1 L2 Memory Disk

<table>
<thead>
<tr>
<th>Type of Storage</th>
<th>L1</th>
<th>L2</th>
<th>Memory</th>
<th>Disk</th>
</tr>
</thead>
<tbody>
<tr>
<td>Typical Size</td>
<td>&lt; 100 KB</td>
<td>&lt; 1 GB</td>
<td>&lt; 10 GB</td>
<td>Many Gb</td>
</tr>
<tr>
<td>Typical Access Time (ns)</td>
<td>25 - 50</td>
<td>5 - 250</td>
<td>50 - 250</td>
<td>5,000,000</td>
</tr>
<tr>
<td>Bandwidth (MB/sec)</td>
<td>10,000 - 50,000</td>
<td>5,000 - 20,000</td>
<td>2,500 - 10,000</td>
<td>50 - 800</td>
</tr>
<tr>
<td>Managed by</td>
<td>Compiler</td>
<td>Hardware</td>
<td>OS</td>
<td>OS</td>
</tr>
</tbody>
</table>


Usually there are two L1 cache - one for Instructions and one for Data. You will often see this referred to in data sheets as “L1 cache: 64KB + 64KB” or “I and D cache.”
Cache Hits and Misses
When the CPU asks for a value from memory, and that value is already in the cache, it can get it quickly. This is called a cache hit.
When the CPU asks for a value from memory, and that value is not already in the cache, it will have to go off the chip to get it. This is called a cache miss.

Possible Cache Architectures
1. Fully Associative – cache lines from any block of memory can appear anywhere in cache.
2. Direct Mapped – a cache line from a particular block of memory has only one place it could appear in cache. A memory block’s cache line is:
   \[ \text{Cache line #} = \text{Memory block #} \mod \text{# lines the cache has} \]
3. N-way Set Associative – a cache line from a particular block of memory can appear in a limited number of places in cache. Each “limited place” is called a set of cache lines. A set contains N cache lines. A memory block’s set number is:
   \[ \text{Set #} = \text{Memory block #} \mod \text{# sets the cache has} \]
   The memory block can appear in any cache line in its set. N is typically 4 for L1 and 8 or 16 for L2.

Memory is arranged in blocks. The size of a block is typically 64 Kbytes. While cache might be multiple kilo- or megabytes, the byes are transferred in much smaller quantities, each called a cache line. The size of a cache line is typically just 64 bytes.

What Happens When the Cache is Full and a New Piece of Memory Needs to Come In?
1. Random – randomly pick a cache line to remove
2. Least Recently Used (LRU) – remove the cache line which has gone unaccessed the longest
3. Oldest (FIFO, First-In-First-Out) – remove the cache line that has been there the longest

Actual Caches Today are N-way Set Associative
This is like a good-news / bad-news joke:

**Good news:** This keeps one specific block of memory from hogging all the cache lines.

**Bad news:** It also means that a block that your program uses much more than the other blocks will need to over-use those N cache lines assigned to that block, and underuse the other cache lines assigned to the other blocks.

Actual Cache Architecture
Let’s try some reasonable numbers. Assume there is 1 GB of main memory, 512 KB of L2 cache, and 64-byte 16-way cache lines in the L2 cache. This means that there are

\[ \frac{512K}{64} = 8192 \]
cache lines available all together in the cache. The “16-way” means that there are 16 cache lines per set, so that the cache must be divided into:

\[ \frac{8192}{16} = 512 \]
sets, which makes each memory block need to contain:

\[ \frac{1GB}{512} = \frac{1,073,741,824}{512} = 2,097,152 = 2MB \]
of main memory. What does this mean to you?

Actual Cache Architecture – Take Home Message
A 2 MB memory block could be contained in

\[ \frac{2MB}{64} = \frac{524,288}{64} = 32,768 \]
cache lines, but it only has 16 available. If your program is using data from just one block, and they are coming from all over the block and going back and forth, you will have many cache misses and those 16 cache lines will keep being re-loaded and re-loaded. Performance will suffer.

If your program is using data from just one block, and they are coming from just a (16 cache-lines) * (64 bytes/cache-line) = 1024-byte portion of the block, those 16 cache lines will never need to be re-loaded. Your cache hit rate will be 100%. Performance will be much better.

To expect the 1024-situation is somewhat unrealistic. However, if you can arrange it so that most of the time you are tripping through your data in-order, then at least you will not have nearly as many cache misses. It is the jumping around in memory that kills you.

What does this tell you about linked lists?
How Bad Is It? -- Demonstrating the Cache-Miss Problem

C and C++ store 2D arrays a row-at-a-time, like this:

```
0 1 2 3 4
5 6 7 8 9
10 11 12 13 14
15 16 17 18 19
20
```

For large arrays, would it be better to process the elements by row, or by column? Which will avoid the most cache misses?

```c
float f = Array[i][j];
```

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```c
#include <stdio.h>
#include <ctime>
#include <cstdlib>
#define N 10000
float Array[N][N];
double Time();

int main( int argc, char *argv[] )
{
    float sum = 0.;
    double start = Time();
    for( int i = 0; i < N; i++ )
    {
        for( int j = 0; j < N; j++ )
        {
            sum += Array[i][j];  // access across a row
        }
    }
    double finish = Time();
    double row_secs = finish - start;

    sum = 0.;
    start = Time();
    for( int j = 0; j < N; j++ )
    {
        for( int i = 0; i < N; i++ )
        {
            sum += Array[i][j];  // access down a column
        }
    }
    finish = Time();
    double col_secs = finish - start;
    fprintf( stderr, "N = %5d ; By rows = %lf ; By cols = %lf\n", N, row_secs, col_secs );
}

double Time()
{
    return (double)clock() / CLOCKS_PER_SEC;
}
```

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### Time, in seconds, to compute the array sums, based on row-first versus column-first order:

<table>
<thead>
<tr>
<th>N (Thousands)</th>
<th>Row Time</th>
<th>Col Time</th>
</tr>
</thead>
<tbody>
<tr>
<td>10000</td>
<td>0.87</td>
<td>0.36</td>
</tr>
<tr>
<td>15000</td>
<td>1.94</td>
<td>0.84</td>
</tr>
<tr>
<td>20000</td>
<td>3.47</td>
<td>1.49</td>
</tr>
<tr>
<td>25000</td>
<td>5.45</td>
<td>2.32</td>
</tr>
<tr>
<td>30000</td>
<td>7.85</td>
<td>3.34</td>
</tr>
<tr>
<td>35000</td>
<td>10.71</td>
<td>4.56</td>
</tr>
<tr>
<td>40000</td>
<td>14.69</td>
<td>6.59</td>
</tr>
</tbody>
</table>

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Array-of-Structures vs. Structure-of-Arrays:

```
struct xyz { float x, y, z; } Array[N];
```

1. Which is a better use of the cache if we are going to be using X-Y-Z triples a lot?
2. Which is a better use of the cache if we are going to be looking at all X's, then all Y's, then all Z's?

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Sometimes Good Object-Oriented Programming Style is Inconsistent with Good Cache Use:

```
class xyz { public:
            float x, y, z;
            xyz*next;
            xyz();
            static xyz*Head = NULL;
        };
```

It might be better to create a large array of xyz structures and then have the constructor method pull new ones from that list. That would keep many of the elements close together while preserving the flexibility of the linked list.
A function jump table requires two instruction cache lines:
1. The for-loop statements
2. The function that really gets called
and one data cache line:
1. Your funcptr array

Because you create this table, it is in data memory.

What is Brian Apgar’s Complaint About Cache?

The PS2, PSP, and Xbox 360 have 2-way associative caches. Brian talks about a tight loop of indirect function calls being a problem for these systems.
But, those systems have 2-way associative caches for both the instructions and the data. So, there should be enough room.
What, then, does Brian warn us about?

class Bird{
public:
virtual void f();
};
class Raven : public Bird{
public:
void f(); // overrides Bird::f();
};

A Class That Inherits from a Virtual Class uses a Jump Table ("vtable") that is Contained in Instruction Memory

So that the correct f() can be called at runtime, each class has a virtual function table, or vtable, containing a pointer to its own version of f(). Because this vtable is built by the compiler, it ends up living where the compiler wants it, not where you want it.
So, the virtual function jump table requires three instruction cache lines:
1. The for-loop statements
2. The for-loop function that really gets called
3. The virtual function table

The PS2 and PSP have 2-way associative caches. If all of these 3 items are in the same memory block, they will attempt to use the same set in the cache. But, there are only two cache lines in each set. So, every pass through the for-loop will cause a cache miss.
(The Wii has an 8-way associative cache, so there is no similar problem there.)