OpenMP Tasks

Remember OpenMP Sections?

Sections are independent blocks of code, able to be assigned to separate threads if they are available.

```c
#pragma omp parallel sections
{
    #pragma omp section
    { Task 1 }
    #pragma omp section
    { Task 2 }
}
```

There is an implied barrier at the end

OpenMP sections are static, that is, they are good if you know, when you are writing the program, how many of them you will need.

It would be nice to have something more Dynamic

Imagine a capability where you can write something to do down on a Post-It® note, accumulate the Post-It notes, then have all of the threads together execute that set of tasks.

You would also like to not have to know, ahead of time, how many of these Post-It notes you will write. That is, you want the total number to be dynamic.

Well, congratulations, you have just invented OpenMP Tasks!

OpenMP Tasks

- An OpenMP task is a single line of code or a structured block which is immediately "written down" in a list of tasks.
- The new task can be executed immediately, or it can be deferred.
- If the if-clause is used and the argument evaluates to 0, then the task is executed immediately, superseding whatever else that thread is doing.
- There has to be an existing parallel thread team for this to work. Otherwise one thread ends up doing all tasks and you don’t get any contribution to parallelism.
- One of the best uses of this is to process elements of a linked list or a tree.

If You Run This a Number of Times, You Get This:

(What Happened?)

1. Why do we not get the same output every time?
2. Why do we get 4 things printed when we only have print statements in 2 tasks?

Not so simple, huh?

The first answer is easy. Unless you make some special arrangements, the order of execution of the different tasks is undefined.

The second answer is that we actually asked each of the two threads to put two tasks on the sticky notes, for a total of four. How can we get only one thread to do this?
The “single” Pragma

```c
omp_set_num_threads(2);
#pragma omp parallel default(none)
{
    #pragma omp single
    {
        #pragma omp task
        fprintf(stderr, "A\n");
        #pragma omp task
        fprintf(stderr, "B\n");
    }
}
```

When using Tasks, you only want one thread to write the things to do down on the sticky note, but you want all of the threads to be able to execute the sticky notes.

But, if you run this, the order of printing will still be non-deterministic. To solve that problem, do this:

```c
omp_set_num_threads(2);
#pragma omp parallel
{
    #pragma omp single default(none)
    {
        #pragma omp task depend(OUT: a)
        a = 10.;
        #pragma omp task depend(IN: a, OUT: b)
        b = a + 16.;
        #pragma omp task depend(IN: b)
        c = b + 12.;
    }
    #pragma omp taskwait
}
```

Causes all tasks to wait until they are completed.

A Better OpenMP Task Example:

Processing each Element of a Linked List

```c
#pragma omp parallel default(none)
{
    #pragma omp single default(none)
    {
        element *p = listHead;
        while( p != NULL )
        {
            #pragma omp task firstprivate(p)
            Process( p );
            p = p->next;
        }
    }
    #pragma omp taskwait
}
```

Without this, each thread does a full traversal – bad idea!

One more thing – Task Dependencies

```c
omp_set_num_threads(3);
#pragma omp parallel
{
    #pragma omp single default(none)
    {
        float a, b, c;
        #pragma omp task depend(OUT: a)
        a = 10.;
        #pragma omp task depend(IN: a, OUT: b)
        b = a + 16.;
        #pragma omp task depend(IN: b)
        c = b + 12.;
    }
    #pragma omp taskwait
}
```

This maintains the proper dependencies, but, because it involves all of the tasks, it essentially serializes the parallelism out of them. Be careful not to go overboard with dependencies!

Tree Traversal Algorithms

Given a tree:

- We would like to traverse it as quickly as possible.
- We are assuming that we do not need to traverse it in order.
- We just need to visit all nodes.

Given a tree: A

```
  B
 /   \
D     E
  
F     G
  
H I J
  K L M
  N O
```

Given a tree: A

```
  B
 /   \
D     E
  
F     G
  
H I J
  K L M
  N O
```

- This is common in graph algorithms, such as searching.
- If the tree is binary and is balanced, then the maximum depth of the tree is log2(# of Nodes).

Strategies at a node:
1. follow one descendent node
2. follow the other descendent node
3. process the node you’re at

This order could be rearranged, depending on what you are trying to do.
Tree Traversal Algorithms

Without this, thread #0 has to do everything
Without this, each thread does a full traversal – bad idea!

Without this, thread #0 has to do everything
Put this here if you want to wait for all nodes to be traversed before proceeding

Parallelizing a Binary Tree Traversal with Tasks

void Traverse( Node *n )
{
  if( n->left != NULL )
  {
    #pragma omp task private(n) untied
    Traverse( n->left );
  }
  if( n->right != NULL )
  {
    #pragma omp task private(n) untied
    Traverse( n->right );
  }
  #pragma omp taskwait
  Process( n );
}

Parallelizing a Binary Tree Traversal with Tasks:

Tied (gcc 8.2)

How Evenly Tasks Get Assigned to Threads

<table>
<thead>
<tr>
<th>6 Levels – g++ 4.9:</th>
<th>6 Levels – g++ 8.2:</th>
</tr>
</thead>
<tbody>
<tr>
<td>Thread #</td>
<td>Number of Tasks</td>
</tr>
<tr>
<td>0</td>
<td>1</td>
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<tr>
<td>1</td>
<td>32</td>
</tr>
<tr>
<td>2</td>
<td>47</td>
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<tr>
<td>3</td>
<td>47</td>
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<th>12 Levels – g++ 4.9:</th>
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<tr>
<td>Thread #</td>
<td>Number of Tasks</td>
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<td>2813</td>
</tr>
<tr>
<td>3</td>
<td>2815</td>
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Parallelizing a Binary Tree Traversal with Tasks:

Untied (gcc 8.2)
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<td>0</td>
<td>1</td>
</tr>
<tr>
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<td>2048</td>
<td>1</td>
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<tr>
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</tr>
<tr>
<td>3</td>
<td>3071</td>
<td>3</td>
<td>2089</td>
</tr>
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### Benchmarking a Binary Task-driven Tree Traversal

```c
void Process( Node *n )
{
    for( int i = 0; i < 1024; i++ )
    {
        n->value = pow( n->value, 1.1 );
    }
}
```

### Performance vs. Number of Threads

<table>
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<tr>
<th># Threads</th>
<th>Nodes Processed per Second</th>
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### Parallelizing a Tree Traversal with Tasks

- Tasks get spread among the current “thread team”
- Tasks can execute immediately or can be deferred. They are executed at “some time”.
- Tasks can be moved between threads, that is, if one thread has a backlog of tasks to do, an idle thread can come steal some workload.
- Tasks are more dynamic than sections. The task paradigm would still work if there was a variable number of children at each node.

### Parallelizing an N-Tree Traversal with Tasks

```c
void Traverse( Node *n )
{
    for( int i = 0; i < n->numChildren; i++ )
    {
        if( n->child[i] != NULL )
        {
            #pragma omp task
            Traverse( n->child[i] );
        }
    }
    #pragma omp taskwait
    Process( n );
}
```
Performance vs. Number of Levels

Performance vs. Number of Threads

8-thread Speed-up ≈ 6.7
F₀ = ??%
Max Speed-up = ??

\[ \text{Speedup} = \frac{n}{(n-1)} \left( 1 - \frac{1}{\text{Speedup}} \right) = 97\% \]

\[ \text{max Speedup} = \frac{1}{1 - F₀} = 33x \]