A TASTE OF HASKELL

Simon Peyton Jones
Microsoft Research

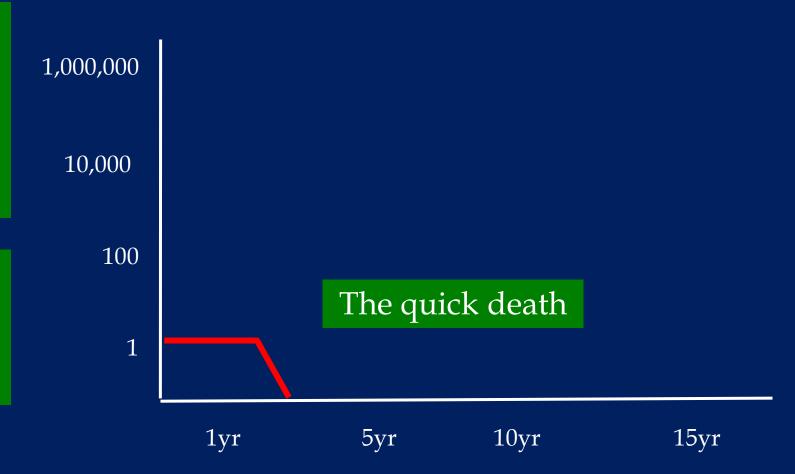
What is Haskell?

- Haskell is a programming language that is
 - purely functional
 - lazy
 - higher order
 - strongly typed
 - general purpose

Why should I care?

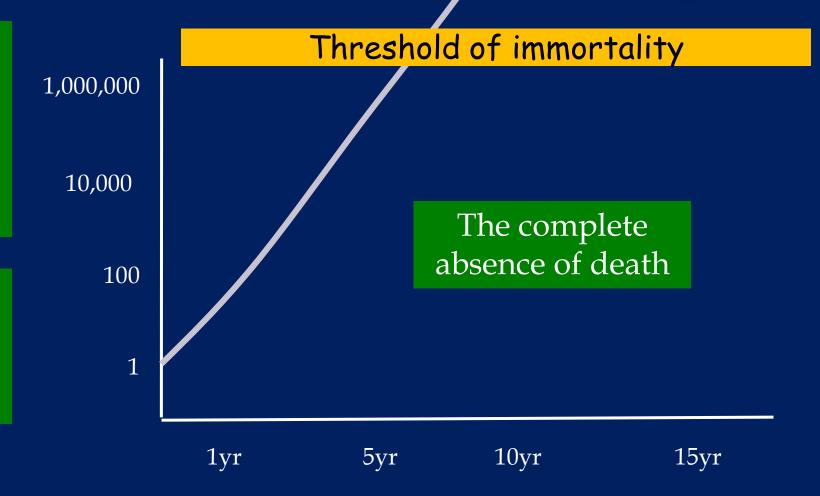
- Functional programming will make you think differently about programming
 - Mainstream languages are all about state
 - Functional programming is all about values
- Whether or not you drink the Haskell Kool-Aid, you'll be a better programmer in whatever language you regularly use

Most research languages





C++, Java, Perl, Ruby



1,000,000

10,000

100

1

1990

Haskell



1995

"Learning Haskell is a great way of training yourself to think functionally so you are ready to take full advantage of C# 3.0 when it comes out" (blog Apr 2007)



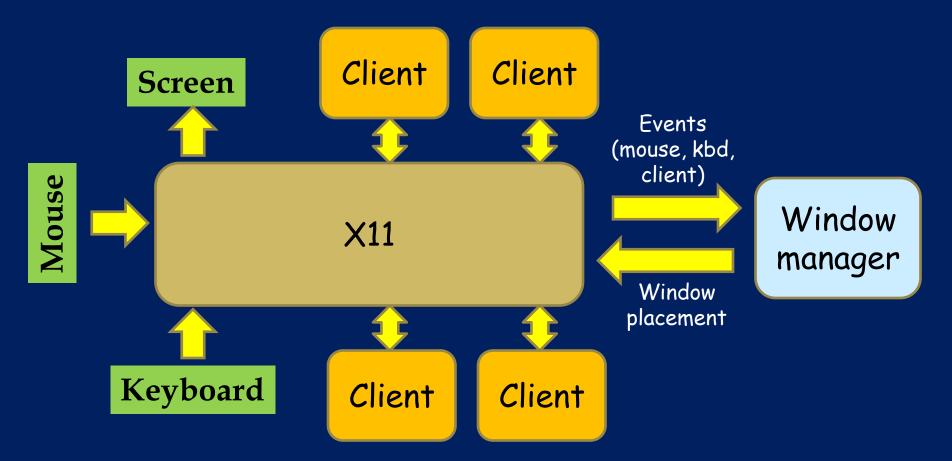
2005

2010

2000

xmonad

xmonad is an X11 tiling window manager written entirely in Haskell



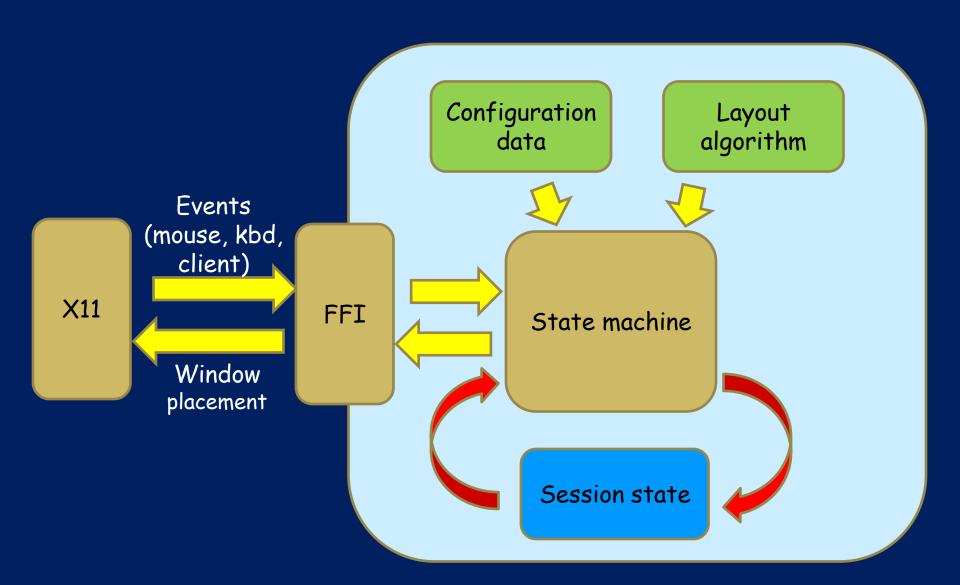
Why I'm using xmonad

- Because it's
 - A real program
 - of manageable size
 - that illustrates many Haskell programming techniques
 - is open-source software
 - is being actively developed
 - by an active community

"Manageable size"

	Code	Comments	Language
metacity	>50k		С
ion3	20k	7k	С
larswm	6k	1.3k	С
wmii	6k	1k	С
dwm 4.2	1.5k	0.2k	С
xmonad 0.2	0.5k	0.7k	Haskell

Inside xmonad



The window stack

Define

new types





Comments

```
module Stack( Stack, insert, swap, ...) where
                                        Import things
import Graphics.X11( Window )
                                      defined elsewhere
type Stack = ...
                                         Specify type
insert :: Window -> Stack -
                                           of insert
-- Newly inserted window has focus
insert = ...
swap :: Stack -> Stack
-- Swap focus with next
swap = \dots
```

The window stack

No import

any more

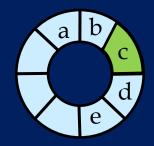
Stack should not exploit the fact that it's a stack of windows

```
module Stack (Stack, insert, swap, ...) where
                                    A stack of values of
type Stack w = ...
                                         type w
insert :: w -> Stack w
-- Newly inserted window has focus
insert = ...
                                      Insert a 'w'
swap :: Stack w -> Stack w
                                      into a stack
-- Swap focus with next
                                        of w's
swap = \dots
```

The window stack

A list takes one of two forms:

- [], the empty list
- (w:ws), a list whose head is w, and tail is ws



A ring of windows
One has the focus

The type "[w]" means "list of w"

The ring above is represented [c,d,e,...,a,b]

Functions are defined by pattern matching

w1:w2:ws means w1:(w2:ws)

Syntactic sugar

Running Haskell

- Download:
 - ghc: http://haskell.org/ghc
 - Hugs: http://haskell.org/hugs
- Interactive:
 - ghci Stack.hs
 - hugs Stack.hs
- Compiled:
 - ghc -c Stack.hs

Rotating the windows

A ring of windows
One has the focus

Pattern matching forces us to think of all cases

```
focusNext :: Stack -> Stack
focusNext (w:ws) = ws ++ [w]
focusnext [] = []
```

Type says "this function takes two arguments, of type [a], and returns a result of type [a]"

```
(++) :: [a] -> [a] -> [a]
-- List append; e.g. [1,2] ++ [4,5] = [1,2,4,5]
```

Definition in Prelude (implicitly imported)

Recursion

```
(++) :: [a] -> [a] -> [a]

-- List append; e.g. [1,2] ++ [,5] = [1,2,4,5]

[] ++ ys = ys

(x:xs) ++ ys = x : (xs ++ ys)
```

Execution model is simple rewriting:

```
[1,2] ++ [4,5]

= (1:2:[]) ++ (4:5:[])

= 1 : ((2:[]) ++ (4:5:[]))

= 1 : 2 : ([] ++ (4:5:[]))

= 1 : 2 : 4 : 5 : []
```

Rotating backwards



```
focusPrev :: Stack -> Stack
focusPrev ws = reverse (focusNext (reverse ws))
```

```
reverse :: [a] -> [a]
-- e.g. reverse [1,2,3] = [3,2,1]
reverse [] = []
reverse (x:xs) = reverse xs ++ [x]
```

Function application by mere juxtaposition

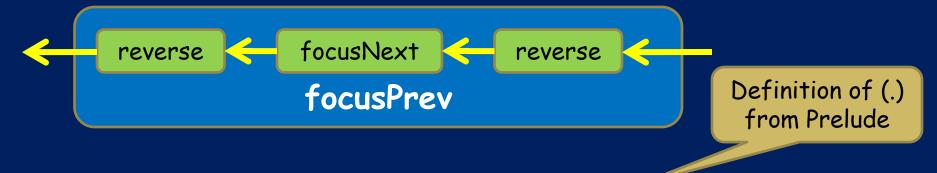
Function application binds more tightly than anything else: (reverse xs) ++ [x]

Function composition

```
focusPrev :: Stack -> Stack
focusPrev ws = reverse (focusNext (reverse ws))
```

can also be written

```
focusPrev :: Stack -> Stack
focusPrev = reverse . focusNext . reverse
```



$$(f . g) x = f (g x)$$

Function composition

Functions as arguments

(.) ::
$$(b->c) -> (a->b) -> (a->c)$$

(f . g) $x = f (g x)$

Just testing

Just testing

- It's good to write tests as you write code
- E.g. focusPrev undoes focusNext; swap undoes itself; etc

```
module Stack where
...definitions...
-- Write properties in Haskell
type TS = Stack Int -- Test at this type

prop_focusNP :: TS -> Bool
prop_focusNP s = focusNext (focusPrev s) == s

prop_swap :: TS -> Bool
prop_swap s = swap (swap s) == s
```

Test interactively

Test.QuickCheck is simply a Haskell library (not a "tool")

```
bash$ ghci Stack.hs
Prelude> :m +Test.QuickCheck

Prelude Test.QuickCheck> quickCheck prop_swap
+++ OK, passed 100 tests

Prelude Test.QuickCheck> quickCheck prop_focusNP
+++ OK, passed 100 tests
```

...with a strangelooking type

```
Prelude Test.QuickCheck> :t quickCheck
quickCheck :: Testable prop => prop -> IO ()
```

A 25-line Haskell script

Test batch-mode

runHaskell Foo.hs <args> runs Foo.hs, passing it <args>

Look for "prop_" tests in here

```
bash$ runhaskell QC.hs Stack.hs
prop_swap: +++ OK, passed 100 tests
prop_focusNP: +++ OK, passed 100 tests
```

No side effects. At all.

```
swap :: Stack w -> Stack w
```

 A call to swap returns a new stack; the old one is unaffected.

```
prop_swap s = swap (swap s) == s
```

A variable 's' stands for an immutable value, not for a location whose value can change with time. Think spreadsheets!

Purity makes the interface explicit

```
swap :: Stack w -> Stack w -- Haskell
```

Takes a stack, and returns a stack; that's all

```
void swap( stack s ) /* C */
```

Takes a stack; may modify it; may modify other persistent state; may do I/O

Pure functions are easy to test

```
prop_swap s = swap (swap s) == s
```

- In an imperative or OO language, you have to
 - set up the state of the object, and the external state it reads or writes
 - make the call
 - inspect the state of the object, and the external state
 - perhaps copy part of the object or global state, so that you can use it in the postcondition

Types are everywhere

```
swap :: Stack w -> Stack w
```

- Usual static-typing rant omitted...
- In Haskell, types express high-level design, in the same way that UML diagrams do; with the advantage that the type signatures are machine-checked
- Types are (almost always) optional: type inference fills them in if you leave them out

Improving the design

Improving the design



```
type Stack w = [w]
-- Focus is head of list

enumerate:: Stack w -> [w]
-- Enumerate the windows in layout order
enumerate s = s
```

Changing focus moves the windows around: confusing!

Improving the design A sequence of windows One has the focus

Want: a fixed layout, still with one window having focus

left

Data type declaration

Constructor of the type

Represented as MkStk [b,a] [c,d,e,f,g]

right

```
data Stack w = MkStk [w] [w] -- left and right resp
-- Focus is head of 'right' list
-- Left list is *reversed*
-- INVARIANT: if 'right' is empty, so is 'left'
```

right

Improving the design

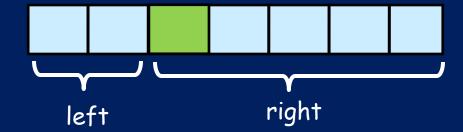
Want: a fixed layout, still with one window having focus

left

Represented as MkStk [b,a] [c,d,e,f,g]

```
data Stack w = MkStk [w] [w] -- left and right resp
-- Focus is head of 'right' list
-- Left list is *reversed*
-- INVARIANT: if 'right' is empty, so is 'left'
enumerate :: Stack w -> [w]
enumerate (MkStack ls rs) = reverse ls ++ rs
```

Moving focus



```
data Stack w = MkStk [w] [w] -- left and right resp

focusPrev :: Stack w -> Stack w

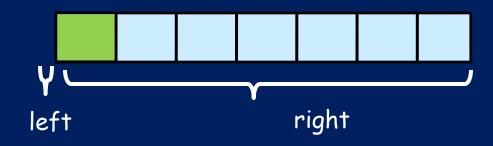
focusPrev (MkStk (l:ls) rs) = MkStk ls (l:rs)

focusPrev (MkStk [] rs) = ...???...
```

Nested pattern matching

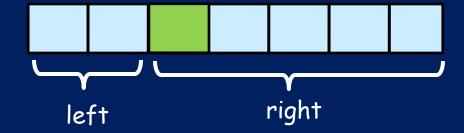
Choices for left=[]:

- no-op
- move focus to end



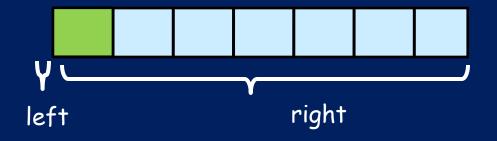
We choose this one

Moving focus



```
data Stack w = MkStk [w] [w] -- left and right resp
-- Focus is head of 'right'

focusPrev :: Stack w -> Stack w
focusPrev (MkStk (1:1s) rs) = MkStk ls (1:rs)
focusPrev (MkStk [] (r:rs)) = MkStk (reverse rs) [r]
```



Choices:

no-op

left

move focus to end



Oops..

```
Warning: Pattern match(es) are non-exhaustive
   In the definition of `focusPrev':
   Patterns not matched: MkStk [] []
```

```
data Stack w = MkStk [w] [w] -- left and right resp
-- Focus is head of 'right'

focusPrev :: Stack w -> Stack w
focusPrev (MkStk (l:ls) rs) = MkStk ls (l:rs)
focusPrev (MkStk [] (r:rs)) = MkStk (reverse rs) [r]
focusPrev (MkStk [] []) = MkStk [] []
```

- Pattern matching forces us to confront all the cases
- Efficiency note: reverse costs O(n), but that only happens once every n calls to focusPrev, so amortised cost is O(1).

Data types

- A new data type has one or more constructors
- Each constructor has zero or more arguments

```
data Stack w = MkStk [w] [w]

data Bool = False | True

data Colour = Red | Green | Blue

data Maybe a = Nothing | Just a
```

Built-in syntactic sugar for lists, but otherwise lists are just another data type

Data types

```
data Stack w = MkStk [w] [w]

data Bool = False | True

data Colour = Red | Green | Blue

data Maybe a = Nothing | Just a
```

- Constructors are used:
 - as a function to construct values ("right hand side")
 - in patterns to deconstruct values ("left hand side")

```
isRed :: Colour -> Bool
isRed Red = True
isRed Green = False
isRed Blue = False
Patterns Values
```

Data types

```
data Maybe a = Nothing | Just a

data Stack w = MkStk [w] [w]
-- Invariant for (MkStk ls rs)
-- rs is empty => ls is empty
```

- Data types are used
 - to describe data (obviously)
 - to describe "outcomes" or "control"

```
module Stack( focus, ...) where

focus :: Stack w -> Maybe w
-- Returns the focused window of the stack
-- or Nothing if the stack is empty
focus (MkStk _ []) = Nothing
focus (MkStk _ (w:_)) = Just w
```

A bit like an exception...

...but you can't forget to catch it No "null-pointer dereference" exceptions

```
module Foo where exceptions import Stack

foo s = ...case (focus s) of

Nothing -> ...do this in empty case...

Just w -> ...do this when there is a focus...
```

Data type abstraction



```
module Operations( ... ) where
import Stack( Stack, focusNext )
f :: Stack w -> Stack w
f (MkStk as bs) = ...
```

OK: Stack is imported

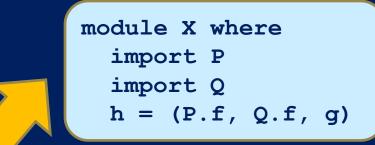
NOT OK: MkStk is not imported

Stack is exported, but not its constructors; so its representation is hidden

Haskell's module system

module P(f,g)

Module system is merely a name-space control mechanism



- Compiler typically does lots of cross-module inlining
- Modules can be grouped into packages



```
module Q(f) where
f = ...
```

```
module Z where
f = ...
```

where

Type classes

The need for type classes

```
delete :: Stack w -> w -> Stack w
-- Remove a window from the stack
```

Can this work for ANY type w?

```
delete :: ∀w. Stack w -> w -> Stack w
```

No - only for w's that support equality

```
sort :: [a] -> [a]
-- Sort the list
```

- Can this work for ANY type a?
- No only for a's that support ordering

The need for type classes

```
serialise :: a -> String
-- Serialise a value into a string
```

Only for w's that support serialisation

```
square :: n -> n
square x = x*x
```

- Only for numbers that support multiplication
- But square should work for any number that does; e.g. Int, Integer, Float, Double, Rational

"for all types w that support the Eq operations"

Type classes

```
delete :: \forall w. Eq w => Stack w -> w -> Stack w
```

If a function works for every type that has particular properties, the type of the function says just that

```
sort :: Ord a => [a] -> [a]
serialise :: Show a => a -> String
square :: Num n => n -> n
```

Otherwise, it must work for any type whatsoever

```
reverse :: [a] -> [a]
filter :: (a -> Bool) -> [a] -> [a]
```

Works for any type 'n' that supports the Num operations

Type classes

```
FORGET all
you know
about OO
classes!
```

```
square :: Num n => n -> n
square x = x*x
```

```
instance Num Int where
a + b = plusInt a b
a * b = mulInt a b
negate a = negInt a
...etc..
```

The class

declaration says

what the Num

operations are

An instance declaration for a type T says how the Num operations are implemented on T's

```
plusInt :: Int -> Int -> Int
mulInt :: Int -> Int -> Int
etc, defined as primitives
```

How type classes work

When you write this...

```
square :: Num n => n -> n
square x = x*x
```

...the compiler generates this

```
square :: Num n \rightarrow n \rightarrow n square d x = (*) d x x
```

The "Num n =>" turns into an extra value argument to the function.

It is a value of data type Num n

A value of type (Num T) is a vector of the Num operations for type T

How type classes work

When you write this...

```
square :: Num n => n -> n
square x = x*x
```

The class decl translates to:

- A data type decl for Num
- A selector function for each class operation

...the compiler generates this

```
square :: Num n \rightarrow n \rightarrow n square d x = (*) d x x
```

A value of type (Num T) is a vector of the Num operations for type T

How type classes work

When you write this...

```
square :: Num n => n -> n
square x = x*x
```

instance Num Int where a + b = plusInt a b a * b = mulInt a b

```
negate a = negInt a
```

```
...etc..
```

...the compiler generates this

```
square :: Num n \rightarrow n \rightarrow n square d x = (*) d x x
```

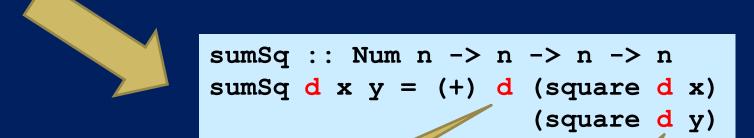
An instance decl for type T translates to a value declaration for the Num dictionary for T

A value of type (Num T) is a vector of the Num operations for type T

All this scales up nicely

 You can build big overloaded functions by calling smaller overloaded functions

```
sumSq :: Num n => n -> n -> n
sumSq x y = square x + square y
```



Extract addition operation from d

Pass on d to square

All this scales up nicely

 You can build big instances by building on smaller instances

Example: complex numbers

```
class Num a where
  (+) :: a -> a -> a
  (-) :: a -> a -> a
  fromInteger :: Integer -> a
  ....

inc :: Num a => a -> a
inc x = x + 1
```

Even literals are overloaded

"1" means
"fromInteger 1"

```
data Cpx a = Cpx a a

instance Num a => Num (Cpx a) where
  (Cpx r1 i1) + (Cpx r2 i2) = Cpx (r1+r2) (i1+i2)
  fromInteger n = Cpx (fromInteger n) 0
```

```
quickCheck :: Test a => a -> IO ()
class Testable a where
  test :: a -> RandSupply -> Bool
class Arbitrary a where
  arby :: RandSupply -> a
instance Testable Bool where
 test b r = b
instance (Arbitrary a, Testable b)
      => Testable (a->b) where
  test f r = test (f (arby r1)) r2
           where (r1,r2) = split r
```

```
prop_swap :: TS -> Bool
```

```
test prop_swap r

= test (prop_swap (arby r1)) r2
where (r1,r2) = split r

Using instance for (->)

Using instance for (->)

Using instance for (->)
```

```
class Arbitrary a where
  arby :: RandSupply -> a
instance Arbitrary Int where
  arby r = randInt r
instance Arbitrary a
                                          Generate Nil value
      => Arbitrary [a] where
  arby r \mid even r1 = []
         | otherwise = arby r2 : arby r3
   where
      (r1,r') = split r
                                           Generate cons value
      (r2,r3) = split r'
```

```
split :: RandSupply -> (RandSupply, RandSupply)
randInt :: RandSupply -> Int
```

- QuickCheck uses type classes to auto-generate
 - random values
 - testing functions

based on the type of the function under test

- Nothing is built into Haskell; QuickCheck is just a library
- Plenty of wrinkles, esp
 - test data should satisfy preconditions
 - generating test data in sparse domains

Type classes = OOP?

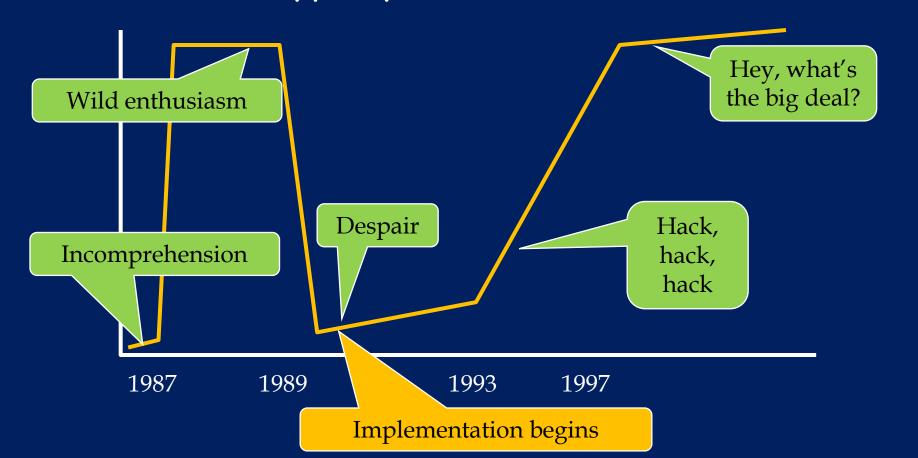
- In OOP, a value carries a method suite
- With type classes, the method suite travels separately from the value
 - Old types can be made instances of new type classes (e.g. introduce new Serialise class, make existing types an instance of it)
 - Method suite can depend on result type e.g. fromInteger :: Num a => Integer -> a
 - Polymorphism, not subtyping

Type classes have proved extraordinarily convenient in practice

- Equality, ordering, serialisation
- Numerical operations. Even numeric constants are overloaded; e.g. f x = x*2
- And on and on....time-varying values, pretty-printing, collections, reflection, generic programming, marshalling, monads, monad transformers....

Type classes over time

 Type classes are the most unusual feature of Haskell's type system



Type-class fertility

Higher kinded type variables (1995)

Implicit parameters (2000)

Extensible

records (1996)

Wadler/ Blott type classes (1989)

Multiparameter type classes (1991)

Functional dependencies (2000)

Computation at the type level

Overlapping instances

> "newtype deriving"

Derivable

Generic programming

Associated types (2005) Testing

Applications

type classes Variations

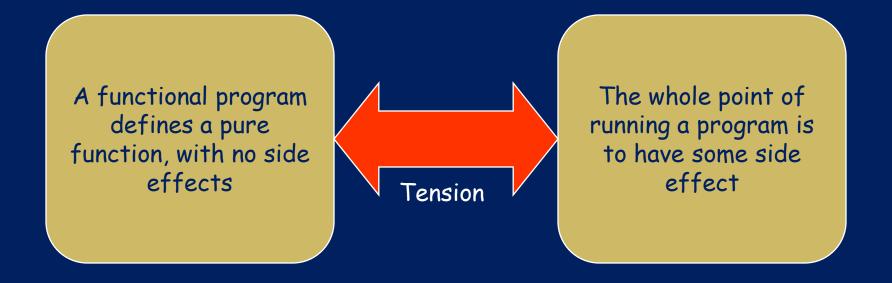
Type classes summary

- A much more far-reaching idea than we first realised: the automatic, type-driven generation of executable "evidence"
- Many interesting generalisations, still being explored
- Variants adopted in Isabel, Clean, Mercury, Hal, Escher
- Long term impact yet to become clear

Doing I/O

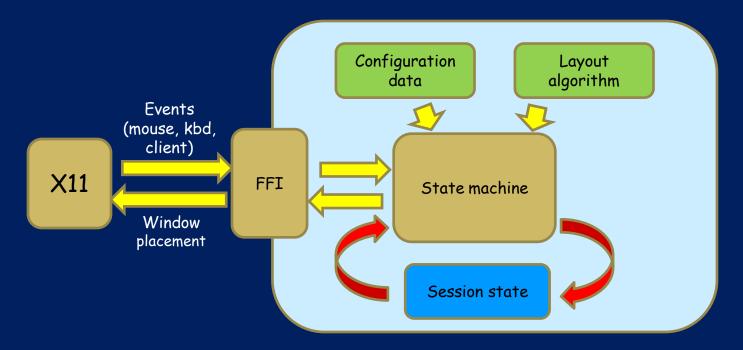
Where is the I/O in xmonad?

- All this pure stuff is very well, but sooner or later we have to
 - talk to X11, whose interface is not at all pure
 - do input/output (other programs)



Where is the I/O in xmonad?

- All this pure stuff is very well, but sooner or later we have to
 - talk to X11, whose interface is not at all pure
 - do input/output (other programs)



Doing I/O

- putStr :: String -> ()
 -- Print a string on the console
- BUT: now swap :: Stack w -> Stack w might do arbitrary stateful things



And what does this do?

```
[putStr "yes", putStr "no"]
```

- What order are the things printed?
- Are they printed at all?

Order of evaluation!

Laziness!

The main idea

A value of type (IO t) is an "action" that, when performed, may do some input/output before delivering a result of type t.

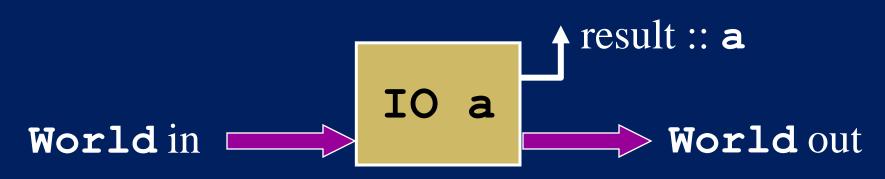
```
putStr :: String -> IO ()
-- Print a string on the console
```

- "Actions" sometimes called "computations"
- An action is a first class value
- Evaluating an action has no effect;
 performing the action has an effect

A helpful picture

A value of type (IO t) is an "action" that, when performed, may do some input/output before delivering a result of type t.

```
type IO a = World -> (a, World)
-- An approximation
```



Simple I/O



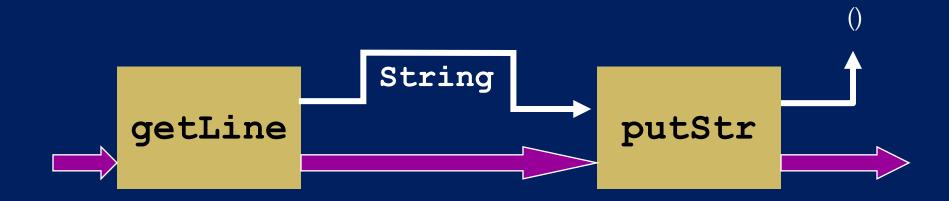
```
getLine :: IO String
```

putStr :: String -> IO ()

Main program is an action of type IO ()

```
main :: IO ()
main = putStr "Hello world"
```

Connecting actions up



Goal: read a line and then write it back out

Connecting actions up

We have connected two actions to make a new, bigger action.

Getting two lines

We want to just return (s1,s2)

The return combinator

```
getTwoLines :: IO (String,String)
getTwoLines = do { s1 <- getLine
    ; s2 <- getLine
    ; return (s1, s2) }</pre>
```

```
return :: a -> IO a
```



Desugaring do notation

- · "do" notation adds only syntactic sugar
- Deliberately imperative look and feel

```
do \{x < -e; s\} = e >>= (\x -> do \{s\})

do \{e\} = e
```

```
(>>=) :: IO a -> (a -> IO b) -> IO b
```

Desugaring "do" notation

```
echo :: IO ()
echo = do { 1 <- getLine; putStr 1 }

echo = getLine >>= (\1 -> putStr 1)
```

A "lambda abstraction" (\times -> e) means "a function taking one parameter, \times , and returning e"

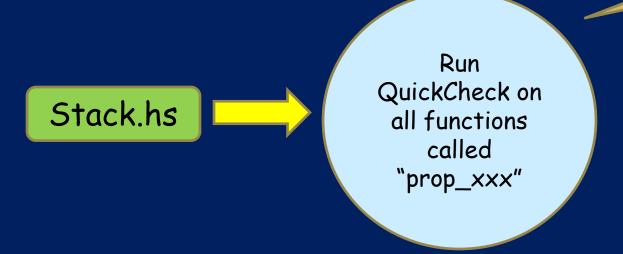
```
(>>=) :: IO a -> (a -> IO b) -> IO b
```

Using layout instead of braces

- You can use
 - explicit braces/semicolons
 - or layout
 - or any mixture of the two

An example: scripting in Haskell

Write this script in Haskell



```
bash$ runhaskell QC.hs Stack.hs
prop_swap: +++ OK, passed 100 tests
prop_focusNP: +++ OK, passed 100 tests
```

```
module Main where
import System; import List
main :: IO ()
main = do { as <- getArgs</pre>
         ; mapM process as }
process :: String -> IO ()
process file = do { cts <- readFile file</pre>
                   ; let tests = getTests cts
                   : if null tests then
                         putStrLn (file ++ ": no properties to check")
                     else do
                   { writeFile "script" $
                        unlines ([":1 " ++ file] ++ concatMap makeTest tests)
                   ; system ("ghci -v0 < script")</pre>
                   ; return () }}
getTests :: String -> [String]
getTests cts = nub $ filter ("prop " `isPrefixOf`) $
                map (fst . head . lex) $ lines cts
makeTest :: String -> [String]
makeTest test = ["putStr \"" ++ p ++ ": \"", "quickCheck " ++ p]
```

Executables have module Main at top

Scripting in Haskell

Import libraries

Module Main must define main :: IO ()

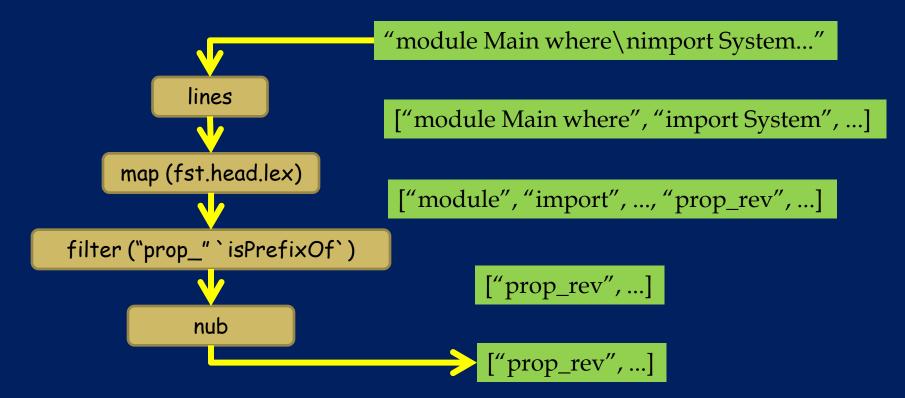
```
getArgs :: IO [String]
-- Gets command line args
```

```
putStrLn :: String -> IO ()
writeFile :: String -> String -> IO ()
system :: String -> IO ExitCode

null :: [a] -> Bool
makeTest :: String -> [String]
concatMap :: (a->[b]) -> [a] -> [b]
unlines :: [String] -> String
```

script

:I Stack.hs putStr "prop_rev" quickCheck prop_rev putStr "prop_focus" quickCheck prop_focus



What have we learned

- Scripting in Haskell is quick and easy (e.g. no need to compile, although you can)
- It is strongly typed; catches many errors
- But there are still many un-handled error conditions (no such file, not lexicallyanalysable, ...)

What have we learned

- Libraries are important; Haskell has a respectable selection
 - Regular expressions
 - Http
 - File-path manipulation
 - Lots of data structures (sets, bags, finite maps etc)
 - GUI toolkits (both bindings to regular toolkits such as Wx and GTK, and more radical approaches)
 - Database bindings

...but not (yet) as many as Perl, Python, C# etc

The types tell the story

type Company = String

I deliver a list of Company

```
sort :: [Company] -> [Company]
-- Sort lexicographically
-- Two calls given the same
-- arguments will give the
-- same results
```

I may do some I/O and then deliver a list of Company

```
sortBySharePrice :: [Company] -> IO [Company]
-- Consult current prices, and sort by them
-- Two calls given the same arguments may not
-- deliver the same results
```

Haskell: the world's finest imperative programming language

- Program divides into a mixture of
 - Purely functional code (most)
 - Necessarily imperative code (some)
- The type system keeps them rigorously separate
- Actions are first class, and that enables new forms of program composition (e.g. mapM_)

First-class control structures

Values of type (IO t) are first class

So we can define our own "control structures"

```
forever :: IO () -> IO ()
forever a = a >> forever a

repeatN :: Int -> IO () -> IO ()
repeatN 0 a = return ()
repeatN n a = a >> repeatN (n-1) a
```

```
e.g. forever (do { e <- getNextEvent ; handleEvent e })
```

Foreign function interface

In the end we have to call C!

This call does not block

Calling convention

Header file and name of C procedure

Haskell

```
foreign import ccall unsafe "HsXlib.h XMapWindow"
  mapWindow :: Display -> Window -> IO ()
```

mapWindow calls XMapWindow

Haskell name and type of imported function

```
void XMapWindow( Display *d, Window *w ) {
   ...
}
```

Marshalling

All the fun is getting data across the border

```
data Display = MkDisplay Addr#
data Window = MkWindow Addr#
```

Addr#: a built-in type representing a C pointer

```
foreign import ccall unsafe "HsXlib.h XMapWindow"
  mapWindow :: Display -> Window -> IO ()
```

'foreign import' knows how to unwrap a single-constructor type, and pass it to C

Marshalling

All the fun is getting data across the border

```
data Display = MkDisplay Addr#
data XEventPtr = MkXEvent Addr#

foreign import ccall safe "HsXlib.h XNextEvent"
    xNextEvent:: Display -> XEventPtr -> IO ()
```

But what we want is

Marshalling

Getting what we want is tedious...

...but there are tools that automate much of the grotesque pain (hsc2hs, c2hs etc).

The rest of Haskell

Laziness

- Haskell is a lazy language

```
cond :: Bool -> a -> a -> a
cond True  t e = t
cond False t e = e
```

Same with local definitions

NB: new syntax guards

Why laziness is important

- Laziness supports modular programming
- Programmer-written functions instead of built-in language constructs

```
(||) :: Bool -> Bool -> Bool
True || x = True
False || x = x
```

Shortcircuiting "or"

Laziness and modularity

```
tails :: String -> [String]
-- All suffixes of s
tails [] = [[]]
tails (x:xs) = (x:xs) : tails xs
```

```
or :: [Bool] -> Bool
-- (or bs) returns True if any of the bs is True
or [] = False
or (b:bs) = b || or bs
```

Why laziness is important

- Typical paradigm:
 - generate all solutions (an enormous tree)
 - walk the tree to find the solution you want

```
nextMove :: Board -> Move
nextMove b = selectMove allMoves
where
   allMoves = allMovesFrom b
```

A gigantic (perhaps infinite) tree of possible moves

Why laziness is important

- Generally, laziness unifies data with control
- Laziness also keeps Haskell pure, which is a Good Thing

Other language features

Advanced types

- Unboxed types
- Multi-parameter type classes
- Functional dependencies
- GADTs
- Implicit parameters
- Existential types
- etc etc

Concurrent Haskell (threads, communication, synchronisation)

Software Transactional Memory (STM)

Template Haskell (meta programming)

Rewrite rules (domain-specific compiler extensions)

Haskell language

Nested Data Parallel Haskell

Monads, monad transformers, and arrows

Generic programming
One program that works
over lots of different
data structures

Haskell's tool ecosystem

Programming environments (emacs, vim,

Visual Studio)

Debugger

Interpreters

(e.g. GHCi, Hugs)

Space and time profiling

Compilers

(e.g. GHC, Jhc, Yhc)

Generators

- parser (cf yacc)
- lexer (cf lex)
- FFI

Coverage testing

Testing (e.g. QuickCheck, Hunit)

Haskell language

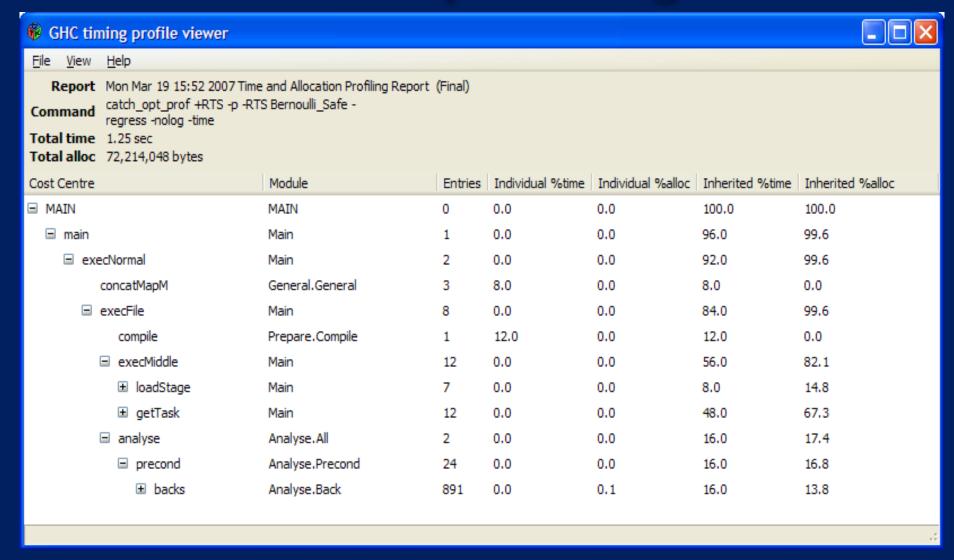
Documentation

generation (Haddock)

Packaging and distribution (Cabal, Hackage)

LIBRARIES

Time profiling



Space profiling

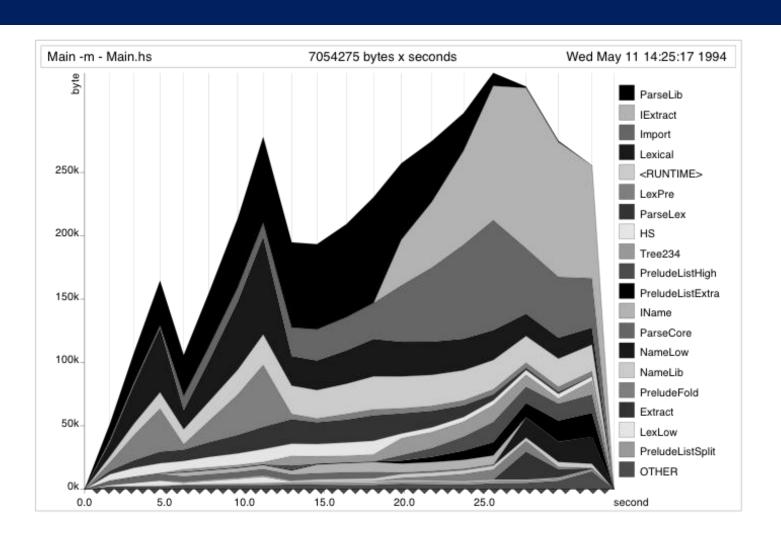
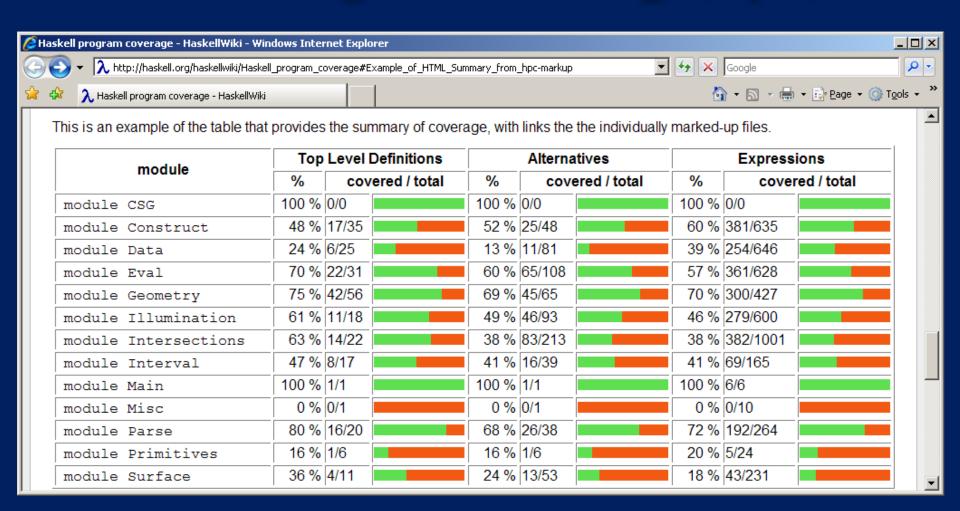


Fig. 18. Heap production of nhe by module, when compiling a small program.

Coverage checking (hpc)

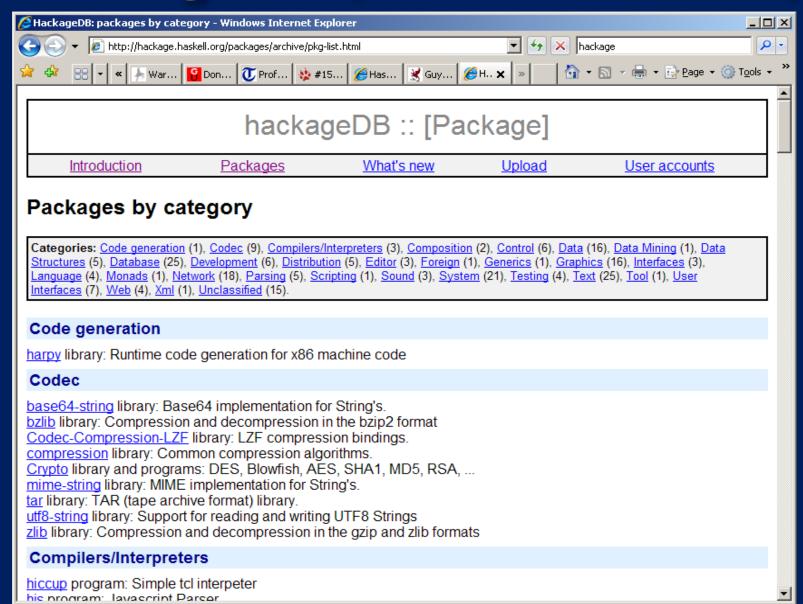


Coverage checking (hpc)

```
reciprocal :: Int -> (String, Int)
                      n > 1 = ('0' : '.' : digits, recur)
2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 8 19 20 1 22 23 24 25 27 28 29 30
    reciprocal n
                      otherwise = error
                         "attempting to compute reciprocal of number <= 1"
      where
       (digits, recur) = divide n 1 []
    divide :: Int -> Int -> [Int] -> (String, Int)
    divide n c cs | c 'elem' cs = ([], position c cs)
                       r == 0 = \frac{(\text{show q, 0})}{r \neq 0} = (\text{show q ++ digits, recur})
      where
     (q, r) = (c*10) \cdot quotRem \cdot n
       (digits, recur) = divide n r (c:cs)
    position :: Int -> [Int] -> Int
    position n (x:xs) | n==x
                            otherwise = 1 + position n xs
    showRecip :: Int -> String
    showRecip n =
       "1/" ++ show n ++ " = " ++
       if r==0 then d else take p d ++ "(" ++ drop p d ++ ")"
      where
      p = length d - r
     (d, r) = reciprocal n
    main = do
                                                    Yellow: not executed
       number <- readLn
                                                    Red: boolean gave False
      putStrLn (showRecip number)
       main
```

Green: boolean gave True

HackageDB (Haskell's CPAN)



Cabal (Haskell's installer)

- A downloaded package, p, comes with
 - p. cabal: a package description
 - Setup.hs: a Haskell script to build/install

```
bash$ ./Setup.hs configure
bash$ ./Setup.hs build
bash$ ./Setup.hs install
```

Standing back...

The central challenge

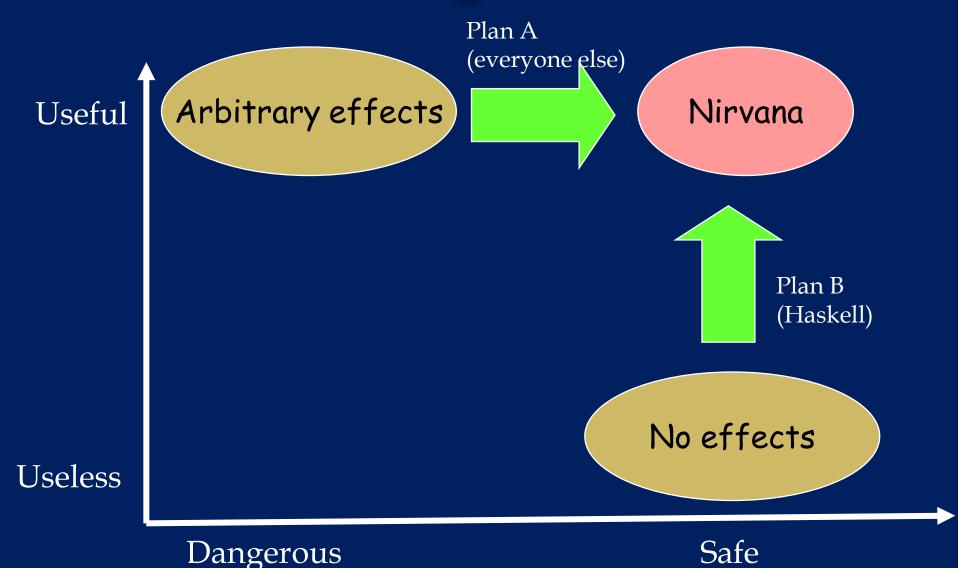
Arbitrary effects Useful Useless

No effects

Dangerous

Safe

The challenge of effects



Two basic approaches: Plan A

Arbitrary effects

Examples

Default = Any effect Plan = Add restrictions

- Regions
- Ownership types
- Vault, Spec#, Cyclone, etc etc

Two basic approaches: Plan B

Default = No effects
Plan = Selectively permit effects

Types play a major role

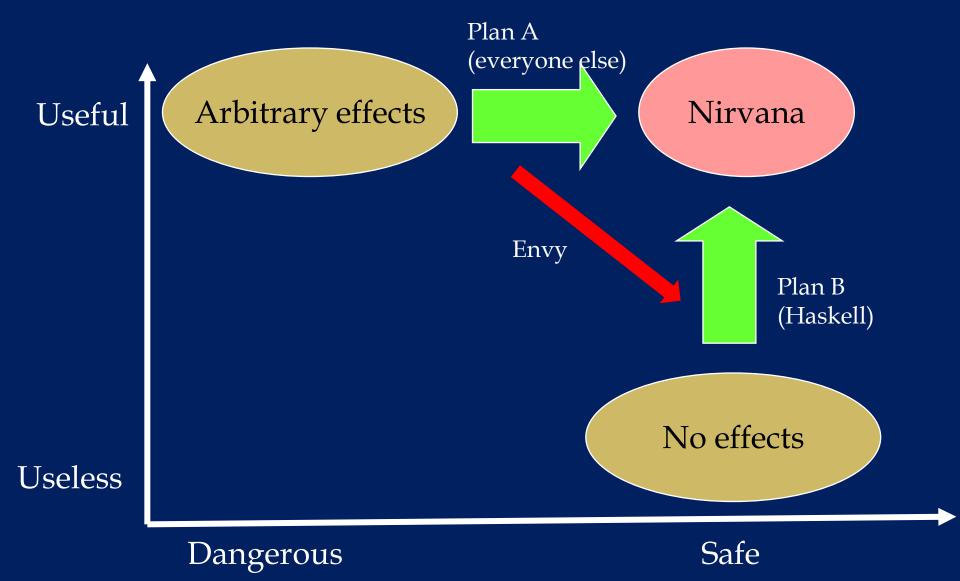
Two main approaches:

- Domain specific languages (SQL, XQuery, MDX, Google map/reduce)
- Wide-spectrum functional languages + controlled effects (e.g. Haskell)

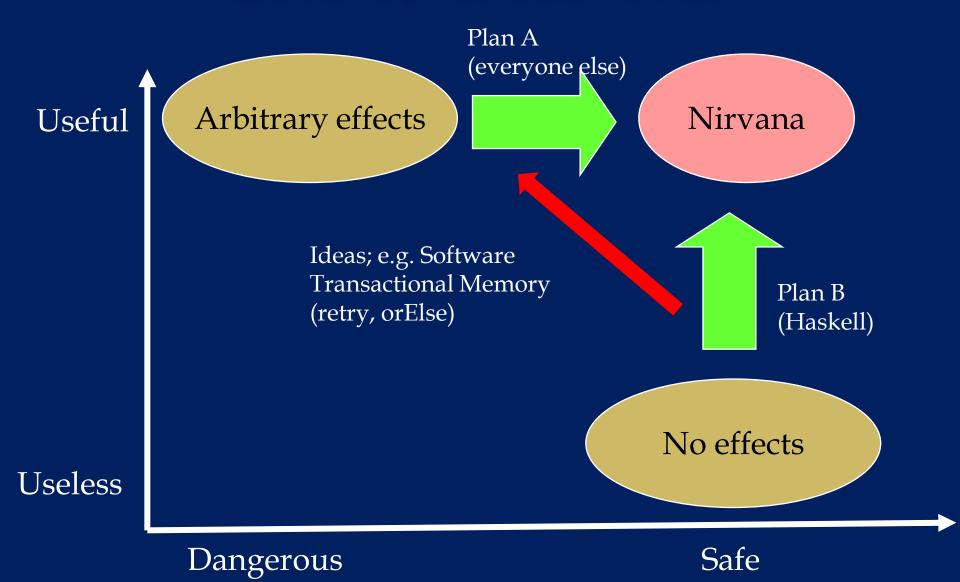


Value oriented programming

Lots of cross-over



Lots of cross-over



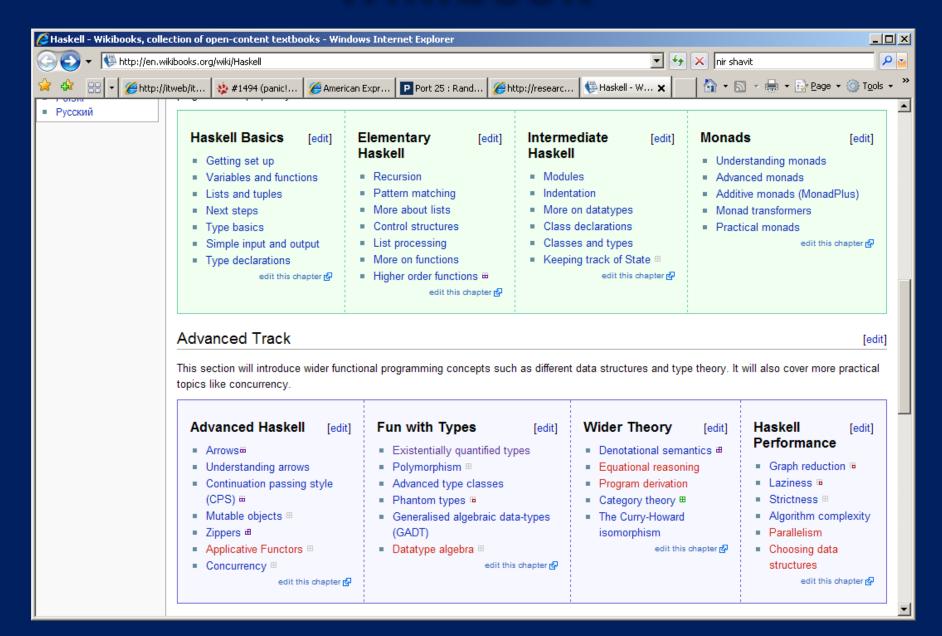
SLPJ conclusions

- One of Haskell's most significant contributions is to take purity seriously, and relentlessly pursue Plan B
- Imperative languages will embody growing (and checkable) pure subsets
- Knowing functional programming makes you a better Java/C#/Perl/Python/Ruby programmer

More info: haskell.org

- The Haskell wikibook
 - http://en.wikibooks.org/wiki/Haskell
- All the Haskell bloggers, sorted by topic
 - http://haskell.org/haskellwiki/Blog_articles
- Collected research papers about Haskell
 - http://haskell.org/haskellwiki/Research_papers
- Wiki articles, by category
 - http://haskell.org/haskellwiki/Category:Haskell
- Books and tutorials
 - http://haskell.org/haskellwiki/Books_and_tutorials

Wikibook



More info: haskell.org

