

References

(“pointers”)

Basic Examples

like C pointers

`r = ref 5`

`malloc`

`!r`

`*r`

`r := 7`

`*r = 7`

`(r:=succ(!r); !r)`

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Basic Examples

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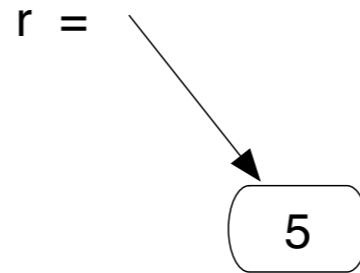
`(r:=succ(!r); r:=succ(!r); r:=succ(!r);
r:=succ(!r); !r)`

i.e.,

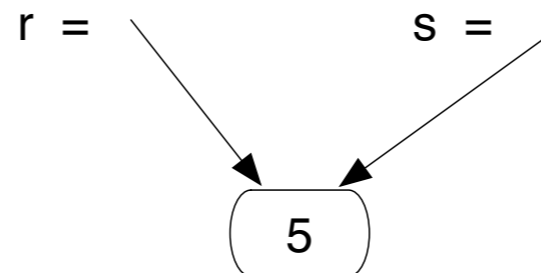
`((((r:=succ(!r); r:=succ(!r)); r:=succ(!r));
r:=succ(!r)); !r)`

Aliasing

A value of type `Ref T` is a *pointer* to a cell holding a value of type `T`.



If this value is “copied” by assigning it to another variable, the cell pointed to is not copied.



So we can change `r` by assigning to `s`:

```
(s:=6; !r)
```

Aliasing all around us

Reference cells are not the only language feature that introduces the possibility of aliasing.

- ▶ arrays
- ▶ communication channels
- ▶ I/O devices (disks, etc.)

The difficulties of aliasing

The possibility of aliasing invalidates all sorts of useful forms of reasoning about programs, both by programmers...

The function

$\lambda r:Ref\ Nat. \lambda s:Ref\ Nat. (r:=2; s:=3; !r)$

always returns 2 unless r and s are aliases.

...and by compilers:

Code motion out of loops, common subexpression elimination, allocation of variables to registers, and detection of uninitialized variables all depend upon the compiler knowing which objects a load or a store operation could reference.

High-performance compilers spend significant energy on *alias analysis* to try to establish when different variables cannot possibly refer to the same storage.

The benefits of aliasing

The problems of aliasing have led some language designers simply to disallow it (e.g., Haskell).

But there are good reasons why most languages do provide constructs involving aliasing:

- ▶ efficiency (e.g., arrays)
- ▶ “action at a distance” (e.g., symbol tables)
- ▶ shared resources (e.g., locks) in concurrent systems
- ▶ etc.

Example

`c = ref 0`

`incc = λx:Unit. (c := succ (!c); !c)` **++c**

`decc = λx:Unit. (c := pred (!c); !c)`

`incc unit`

`decc unit`

`o = {i = incc, d = decc}`

Q: how about c++?


```
let newcounter =  
  λ_:Unit.  
    let c = ref 0 in  
    let incc = λx:Unit. (c := succ (!c); !c) in  
    let decc = λx:Unit. (c := pred (!c); !c) in  
    let o = {i = incc, d = decc} in  
    o
```

Syntax

$t ::=$

unit

x

$\lambda x:T.t$

$t\ t$

$\text{ref } t$

$!t$

$t := t$

terms

unit constant

variable

abstraction

application

reference creation

dereference

assignment

... plus other familiar types, in examples.

Typing Rules

$$\frac{\Gamma \vdash t_1 : T_1}{\Gamma \vdash \text{ref } t_1 : \text{Ref } T_1} \quad (\text{T-REF})$$

$$\frac{\Gamma \vdash t_1 : \text{Ref } T_1}{\Gamma \vdash !t_1 : T_1} \quad (\text{T-DEREF})$$

$$\frac{\Gamma \vdash t_1 : \text{Ref } T_1 \quad \Gamma \vdash t_2 : T_1}{\Gamma \vdash t_1 := t_2 : \text{Unit}} \quad (\text{T-ASSIGN})$$

Final example

```
NatArray = Ref (Nat → Nat);
```

```
newarray = λ_:Unit. ref (λn:Nat.0);  
          : Unit → NatArray
```

```
lookup = λa:NatArray. λn:Nat. (!a) n;  
        : NatArray → Nat → Nat
```

```
update = λa:NatArray. λm:Nat. λv:Nat.  
        let oldf = !a in  
        a := (λn:Nat. if equal m n then v else oldf n);  
        : NatArray → Nat → Nat → Unit
```

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Crucial observation: evaluating `ref 0` must *do* something.

Otherwise,

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r = ref 0
```

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s = ref 0
```

and

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So what is a reference?

The Store

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The Store

A reference names a *location* in the *store* (also known as the *heap* or just the *memory*).

What is the store?

- ▶ *Concretely*: An array of 8-bit bytes, indexed by 32-bit integers.
- ▶ *More abstractly*: an array of *values*
- ▶ *Even more abstractly*: a partial function from *locations* to *values*.

Locations

Syntax of values:

$v ::=$

`unit`

`$\lambda x:T.t$`

`/`

values

unit constant

abstraction value

store location

... and since all values are terms...

Syntax of Terms

$t ::=$

unit

x

$\lambda x:T.t$

$t\ t$

$\text{ref } t$

$!t$

$t := t$

$/$

terms

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store location

Aside

Does this mean we are going to allow programmers to write explicit locations in their programs??

No: This is just a modeling trick. We are enriching the “source language” to include some run-time structures, so that we can continue to formalize evaluation as a relation between source terms.

Aside: If we formalize evaluation in the big-step style, then we can add locations to the set of values (results of evaluation) without adding them to the set of terms.

Evaluation

The result of evaluating a term now depends on the store in which it is evaluated. Moreover, the result of evaluating a term is not just a value — we must also keep track of the changes that get made to the store.

I.e., the evaluation relation should now map a term and a store to a reduced term and a new store.

$$t \mid \mu \longrightarrow t' \mid \mu'$$

We use the metavariable μ to range over stores.

Evaluation

An assignment $t_1 := t_2$ first evaluates t_1 and t_2 until they become values...

$$\frac{t_1 \mid \mu \longrightarrow t'_1 \mid \mu'}{t_1 := t_2 \mid \mu \longrightarrow t'_1 := t_2 \mid \mu'} \quad (\text{E-ASSIGN1})$$

$$\frac{t_2 \mid \mu \longrightarrow t'_2 \mid \mu'}{v_1 := t_2 \mid \mu \longrightarrow v_1 := t'_2 \mid \mu'} \quad (\text{E-ASSIGN2})$$

... and then returns `unit` and updates the store:

$$l := v_2 \mid \mu \longrightarrow \text{unit} \mid [l \mapsto v_2]\mu \quad (\text{E-ASSIGN})$$

A term of the form `ref t1` first evaluates inside `t1` until it becomes a value...

$$\frac{t_1 \mid \mu \longrightarrow t'_1 \mid \mu'}{\text{ref } t_1 \mid \mu \longrightarrow \text{ref } t'_1 \mid \mu'} \quad (\text{E-REF})$$

... and then chooses (allocates) a fresh location `l`, augments the store with a binding from `l` to `v1`, and returns `l`:

$$\frac{l \notin \text{dom}(\mu)}{\text{ref } v_1 \mid \mu \longrightarrow l \mid (\mu, l \mapsto v_1)} \quad (\text{E-REFV})$$

`p = malloc(...)`

A term $!t_1$ first evaluates in t_1 until it becomes a value...

$$\frac{t_1 \mid \mu \longrightarrow t'_1 \mid \mu'}{!t_1 \mid \mu \longrightarrow !t'_1 \mid \mu'} \quad (\text{E-DEREF})$$

... and then looks up this value (which must be a location, if the original term was well typed) and returns its contents in the current store:

$$\frac{\mu(l) = v}{!l \mid \mu \longrightarrow v \mid \mu} \quad (\text{E-DEREFLOC})$$

Evaluation rules for function abstraction and application are augmented with stores, but don't do anything with them directly.

$$\frac{t_1 \mid \mu \longrightarrow t'_1 \mid \mu'}{t_1 \ t_2 \mid \mu \longrightarrow t'_1 \ t_2 \mid \mu'} \quad (\text{E-APP1})$$

$$\frac{t_2 \mid \mu \longrightarrow t'_2 \mid \mu'}{v_1 \ t_2 \mid \mu \longrightarrow v_1 \ t'_2 \mid \mu'} \quad (\text{E-APP2})$$

$$(\lambda x:T_{11}. t_{12}) \ v_2 \mid \mu \longrightarrow [x \mapsto v_2]t_{12} \mid \mu \quad (\text{E-APPABS})$$

Aside: garbage collection

Note that we are not modeling garbage collection — the store just grows without bound.

Aside: pointer arithmetic

We can't do any!

Store Typings

Typing Locations

Q: What is the *type* of a *location*?

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A: It depends on the store!

E.g., in the store $(l_1 \mapsto \text{unit}, l_2 \mapsto \text{unit})$, the term $!l_2$ has type `Unit`.

But in the store $(l_1 \mapsto \text{unit}, l_2 \mapsto \lambda x:\text{Unit}.x)$, the term $!l_2$ has type `Unit → Unit`.

Typing Locations — first try

Roughly:

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$$\frac{\Gamma \vdash \mu(l) : T_1}{\Gamma \vdash l : \text{Ref } T_1}$$

More precisely:

$$\frac{\Gamma \mid \mu \vdash \mu(l) : T_1}{\Gamma \mid \mu \vdash l : \text{Ref } T_1}$$

I.e., typing is now a *four*-place relation (between contexts, *stores*, terms, and types).

Problem

However, this rule is not completely satisfactory. For one thing, it can make typing derivations very large!

E.g., if

$$\begin{aligned} &(\mu = l_1 \mapsto \lambda x:\text{Nat}. \text{999}, \\ & \quad l_2 \mapsto \lambda x:\text{Nat}. !l_1 (!l_1 x), \\ & \quad l_3 \mapsto \lambda x:\text{Nat}. !l_2 (!l_2 x), \\ & \quad l_4 \mapsto \lambda x:\text{Nat}. !l_3 (!l_3 x), \\ & \quad l_5 \mapsto \lambda x:\text{Nat}. !l_4 (!l_4 x)), \end{aligned}$$

then how big is the typing derivation for $!l_5$?

Problem!

But wait... it gets worse. Suppose

$$\begin{aligned} &(\mu = l_1 \mapsto \lambda x:\text{Nat}. !l_2 x, \\ & \quad l_2 \mapsto \lambda x:\text{Nat}. !l_1 x), \end{aligned}$$

Now how big is the typing derivation for $!l_2$?

Store Typings

Observation: The typing rules we have chosen for references guarantee that a given location in the store is *always* used to hold values of the *same* type.

These intended types can be collected into a *store typing* — a partial function from locations to types.

E.g., for

$$\begin{aligned}\mu = (&l_1 \mapsto \lambda x:\text{Nat}. 999, \\ &l_2 \mapsto \lambda x:\text{Nat}. !l_1 (!l_1 x), \\ &l_3 \mapsto \lambda x:\text{Nat}. !l_2 (!l_2 x), \\ &l_4 \mapsto \lambda x:\text{Nat}. !l_3 (!l_3 x), \\ &l_5 \mapsto \lambda x:\text{Nat}. !l_4 (!l_4 x)),\end{aligned}$$

A reasonable store typing would be

$$\begin{aligned}\Sigma = (&l_1 \mapsto \text{Nat} \rightarrow \text{Nat}, \\ &l_2 \mapsto \text{Nat} \rightarrow \text{Nat}, \\ &l_3 \mapsto \text{Nat} \rightarrow \text{Nat}, \\ &l_4 \mapsto \text{Nat} \rightarrow \text{Nat}, \\ &l_5 \mapsto \text{Nat} \rightarrow \text{Nat})\end{aligned}$$

Now, suppose we are given a store typing Σ describing the store μ in which we intend to evaluate some term t . Then we can use Σ to look up the types of locations in t instead of calculating them from the values in μ .

$$\frac{\Sigma(l) = T_1}{\Gamma \mid \Sigma \vdash l : \text{Ref } T_1} \quad (\text{T-LOC})$$

I.e., typing is now a four-place relation between contexts, *store typings*, terms, and types.

Final typing rules

$$\frac{\Sigma(l) = T_1}{\Gamma \mid \Sigma \vdash l : \text{Ref } T_1} \quad (\text{T-LOC})$$

$$\frac{\Gamma \mid \Sigma \vdash t_1 : T_1}{\Gamma \mid \Sigma \vdash \text{ref } t_1 : \text{Ref } T_1} \quad (\text{T-REF})$$

$$\frac{\Gamma \mid \Sigma \vdash t_1 : \text{Ref } T_{11}}{\Gamma \mid \Sigma \vdash !t_1 : T_{11}} \quad (\text{T-DEREF})$$

$$\frac{\Gamma \mid \Sigma \vdash t_1 : \text{Ref } T_{11} \quad \Gamma \mid \Sigma \vdash t_2 : T_{11}}{\Gamma \mid \Sigma \vdash t_1 := t_2 : \text{Unit}} \quad (\text{T-ASSIGN})$$

Q: Where do these store typings come from?

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A: When we first typecheck a program, there will be no explicit locations, so we can use an empty store typing.

So, when a new location is created during evaluation,

$$\frac{l \notin \text{dom}(\mu)}{\text{ref } v_1 \mid \mu \longrightarrow l \mid (\mu, l \mapsto v_1)} \quad (\text{E-REFV})$$

we can observe the type of v_1 and extend the “current store typing” appropriately.

Safety

Preservation

First attempt: just add stores and store typings in the appropriate places.

Theorem (?): If $\Gamma \mid \Sigma \vdash t : T$ and $t \mid \mu \longrightarrow t' \mid \mu'$, then $\Gamma \mid \Sigma \vdash t' : T$.

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Theorem (?): If $\Gamma \mid \Sigma \vdash t : T$ and $t \mid \mu \longrightarrow t' \mid \mu'$, then $\Gamma \mid \Sigma \vdash t' : T$. **Wrong!**

Why is this wrong?

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Why is this wrong?

Because Σ and μ here are not constrained to have anything to do with each other!

(Exercise: Construct an example that breaks this statement of preservation.)

Preservation

A store μ is said to be *well typed* with respect to a typing context Γ and a store typing Σ , written $\Gamma \mid \Sigma \vdash \mu$, if $\text{dom}(\mu) = \text{dom}(\Sigma)$ and $\Gamma \mid \Sigma \vdash \mu(l) : \Sigma(l)$ for every $l \in \text{dom}(\mu)$.

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Next attempt:

Theorem (?): If

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$$t \mid \mu \longrightarrow t' \mid \mu'$$

$$\Gamma \mid \Sigma \vdash \mu$$

then $\Gamma \mid \Sigma \vdash t' : T$.

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then $\Gamma \mid \Sigma \vdash t' : T$.

Still wrong!

What's wrong now?

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$$t \mid \mu \longrightarrow t' \mid \mu'$$

$$\Gamma \mid \Sigma \vdash \mu$$

then $\Gamma \mid \Sigma \vdash t' : T$.

Still wrong!

Creation of a new reference cell...

$$\frac{l \notin dom(\mu)}{\text{ref } v_1 \mid \mu \longrightarrow l \mid (\mu, l \mapsto v_1)} \quad (\text{E-REFV})$$

... breaks the correspondence between the store typing and the store.

Preservation (correct version)

Theorem: If

$$\Gamma \mid \Sigma \vdash t : T$$

$$\Gamma \mid \Sigma \vdash \mu$$

$$t \mid \mu \longrightarrow t' \mid \mu'$$

then, for **some** $\Sigma' \supseteq \Sigma$,

$$\Gamma \mid \Sigma' \vdash t' : T$$

$$\Gamma \mid \Sigma' \vdash \mu'.$$

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Theorem: If

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then, for **some** $\Sigma' \supseteq \Sigma$,

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$$\Gamma \mid \Sigma' \vdash \mu'.$$

Proof: Easy extension of the preservation proof for λ_{\rightarrow} .

Progress

Theorem: Suppose t is a closed, well-typed term (that is, $\emptyset \mid \Sigma \vdash t : T$ for some T and Σ). Then either t is a value or else, for any store μ such that $\emptyset \mid \Sigma \vdash \mu$, there is some term t' and store μ' with $t \mid \mu \longrightarrow t' \mid \mu'$.