

# Language Technology

Fall 2014

## Unit 3: Natural Language Learning

### Unsupervised Learning

(EM, forward-backward, inside-outside)

Liang Huang

[liang.huang.sh@gmail.com](mailto:liang.huang.sh@gmail.com)

# Review of Noisy-Channel Model



Application	Input	Output	$p(i)$	$p(o i)$
Machine Translation	$L_1$ word sequences	$L_2$ word sequences	$p(L_1)$ in a language model	translation model
Optical Character Recognition (OCR)	actual text	text with mistakes	prob of language text	model of OCR errors
Part Of Speech (POS) tagging	POS tag sequences	English words	prob of POS sequences	$p(w t)$
Speech recognition	word sequences	speech signal	prob of word sequences	acoustic model

# Example I: Part-of-Speech Tagging

$$P(t \dots t | w \dots w)$$

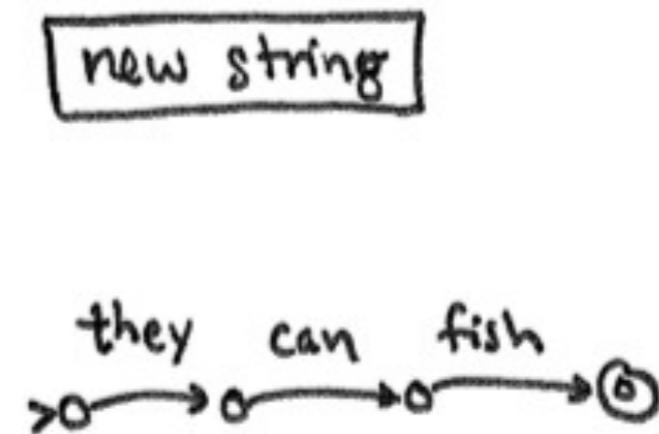
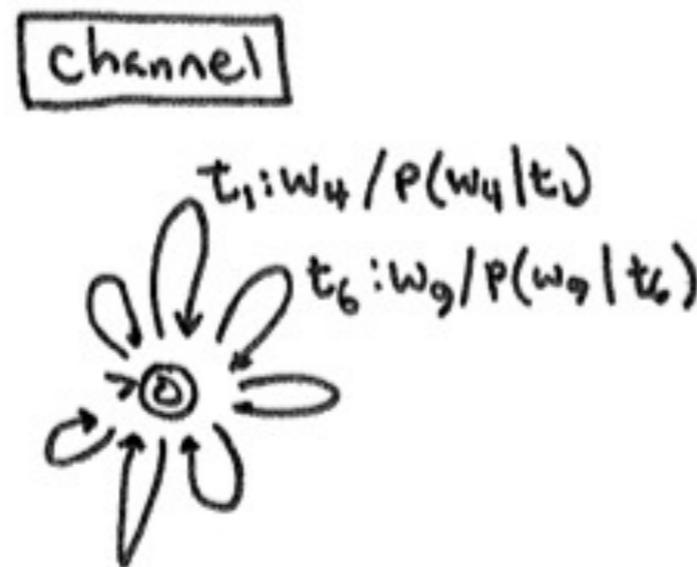
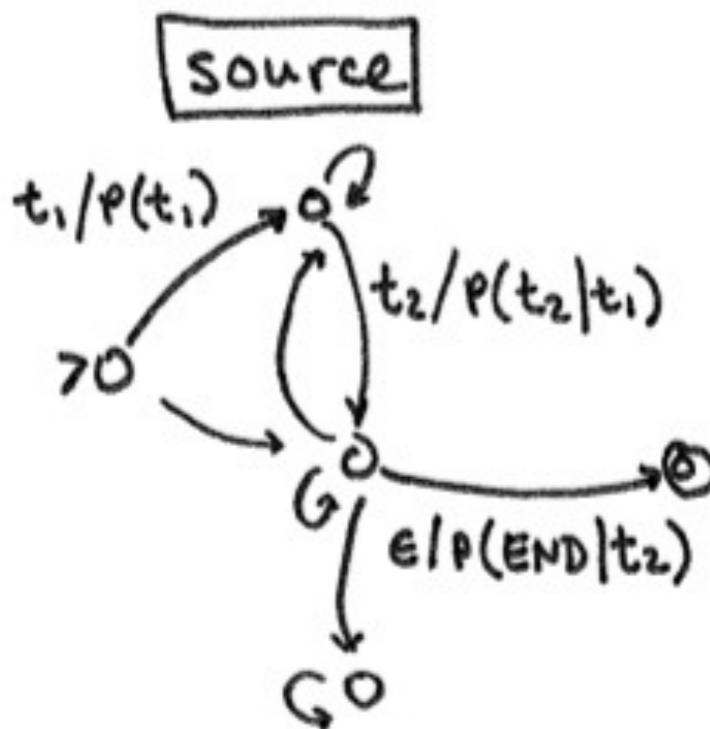
$$\sim P(t \dots t) \cdot P(w \dots w | t \dots t)$$

$$\sim P(t_1) \cdot P(t_2 | t_1) \dots P(t_n | t_{n-1}) \cdot P(w_1 | t_1) \dots P(w_n | t_n)$$

local grammar  
preference

lexical preference

- use tag bigram as a language model
- channel model is context-indep.



# Ideal vs. Available Data



## CRYPTOGRAPHY

1. generate  $e_1, \dots, e_n$  by  $P(e_k | e_{k-1})$
2. for  $i = 1$  to  $n$   
output  $c_i$  by  $P(c_i | e_i)$

ideal

e...

eee...  
|||  
ccc...

available

e...

ccc...

## SPELLING - TO-SOUND

1. generate pho, ... phon
2. transform into  $c_1, \dots, c_m$  by WFST

K A Y E  
/ / ^ \\\  
c a l l e

Y O R A  
^ \ / / !  
l l o r a

K A Y E  
c a l l e

Y O R A  
l l o r a

## MT

1. generate  $e_1, \dots, e_n$  by  $P(e_k | e_{k-1})$
2. for  $i = 1$  to  $n$   
generate  $f_i$  by  $P(f_i | e_i)$
3. permute all  $f_i$  by  $\frac{1}{n!}$

eee...  
| X  
f f f ...

eee...  
X |  
f f f ...

eee...  
f f f ...

eee...  
f f f ...

# Ideal vs. Available Data

## HW2: ideal

EY	B	AH	L	
A	B	E	R	U
1	2	3	4	4

AH	B	AW	T		
A	B	A	U	T	O
1	2	3	3	4	4

AH	L	ER	T		
A	R	A	A	T	O
1	2	3	3	4	4

EY	S		
E	E	S	U
1	1	2	2

## HW4: realistic

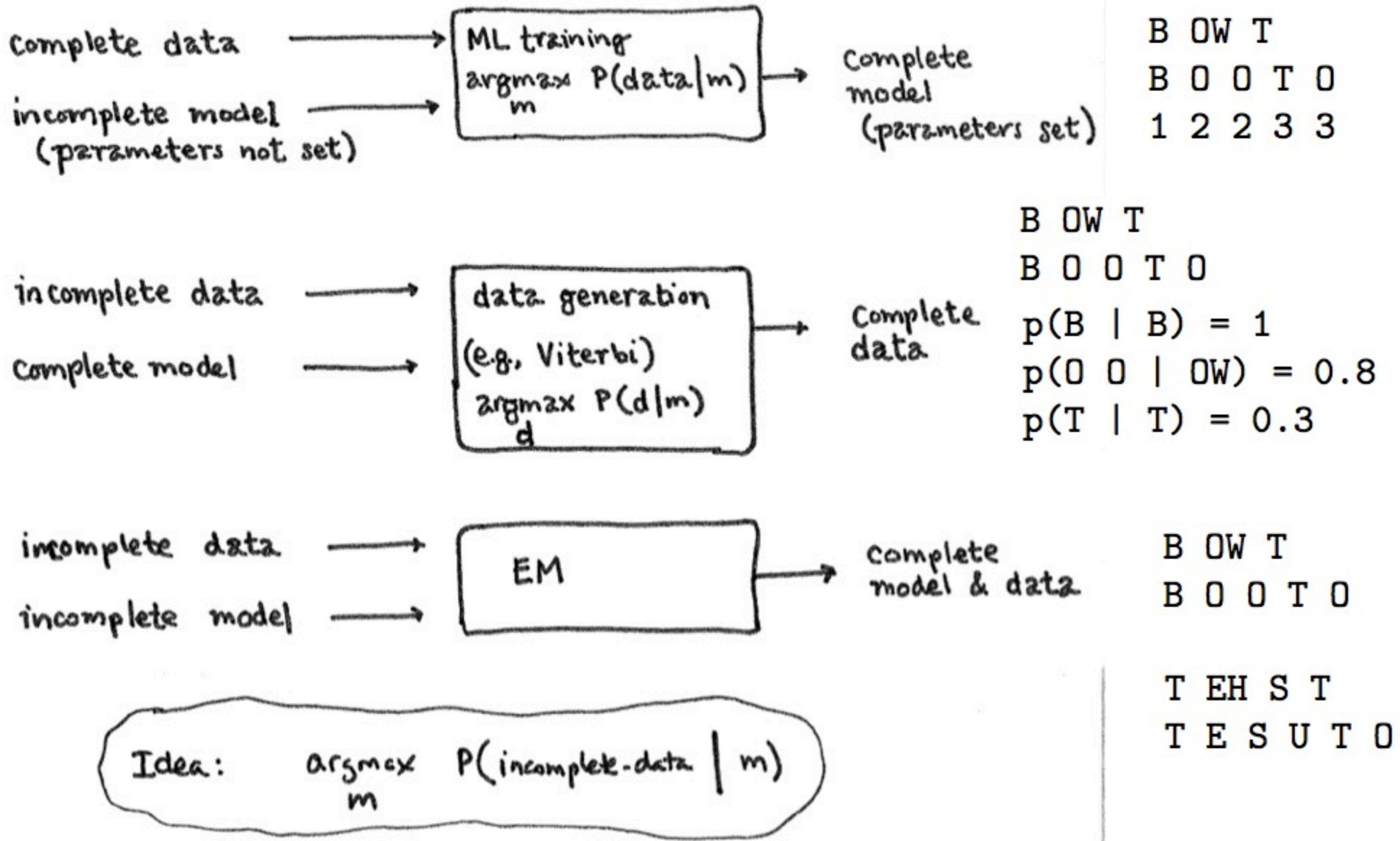
EY	B	AH	L	
A	B	E	R	U

AH	B	AW	T		
A	B	A	U	T	O

AH	L	ER	T		
A	R	A	A	T	O

EY	S		
E	E	S	U

# Incomplete Data / Model



# EM: Expectation-Maximization

Example: Cryptography.  $\underset{m}{\operatorname{argmax}} P(c_1, \dots, c_n | m)$

e...

e...

$$\underset{m}{\operatorname{argmax}} \sum_{e_1, \dots, e_n} P(e_1, \dots, e_n) \cdot P(c_1, \dots, c_n | e_1, \dots, e_n, m)$$

e e e e ...  
i i i i ...  
c c c c ...

ccc ...

$$\underset{m}{\operatorname{argmax}} \sum_{e_1, \dots, e_n} P(e_1, \dots, e_n) \cdot P(c_1 | e_1, m) \cdots P(c_n | e_n, m)$$

each choice of  $m$  yields a specific number!  
some  $m$  are better than others!

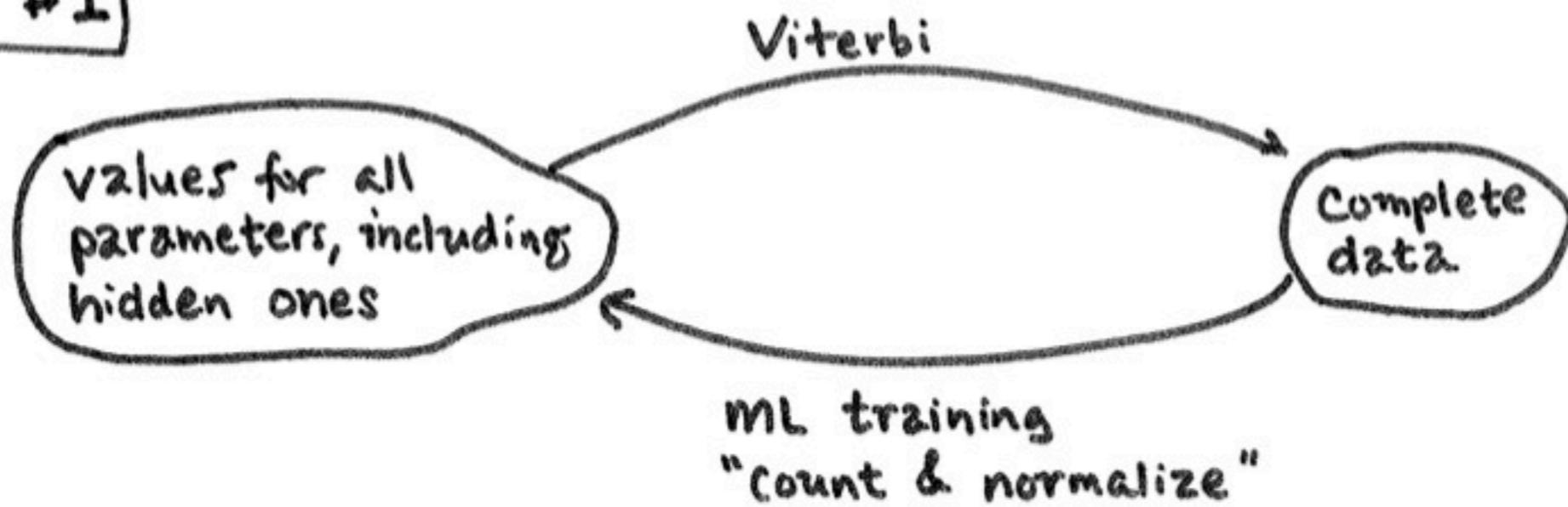
which is best?

EM {

- start with  $m$  such that  $P(c_i | e_j, m) = 1/k$ .
- that gives a certain  $P(c_1, \dots, c_n | m)$ .
- now, change  $m$  to  $m'$  such that  
 $P(c_1, \dots, c_n | m') \geq P(c_1, \dots, c_n | m)$
- (& repeat)

# How to Change $m$ ? I) Hard

Idea #1



Suggests iterative procedure.

initially:

$$\begin{aligned} t(a|x) &= 0.5 \\ t(b|x) &= 0.5 \\ t(a|y) &= 0.5 \\ t(b|y) &= 0.5 \end{aligned}$$

viterbi:    a a a b a a b a a

} NOTE: other decodings are equally good.  
(tie break)

# How to Change m? I) Hard

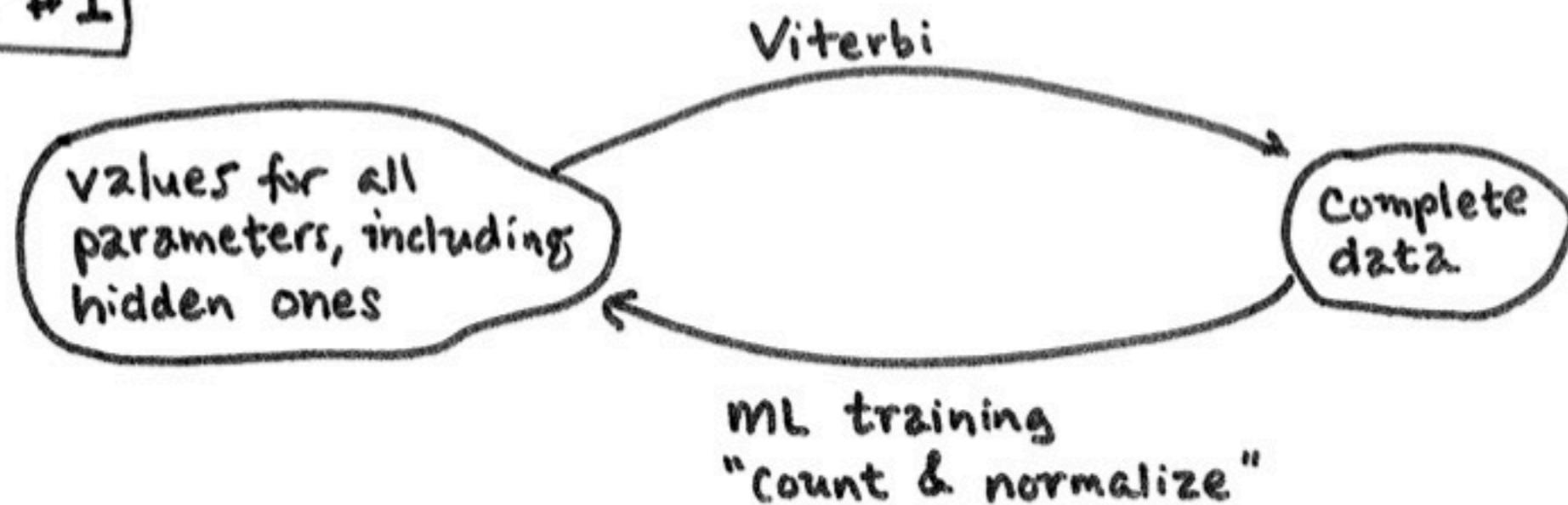
viterbi:      a a a b a a b a a  
                  y x x x x x y x x } NOTE: other  
                                         decodings are  
    equally good.  
    (tie break)

revised:       $t(a|x) = 6/7$   
 $t(b|x) = 1/7$   
 $t(a|y) = 1/2$   
 $t(b|y) = 1/2$

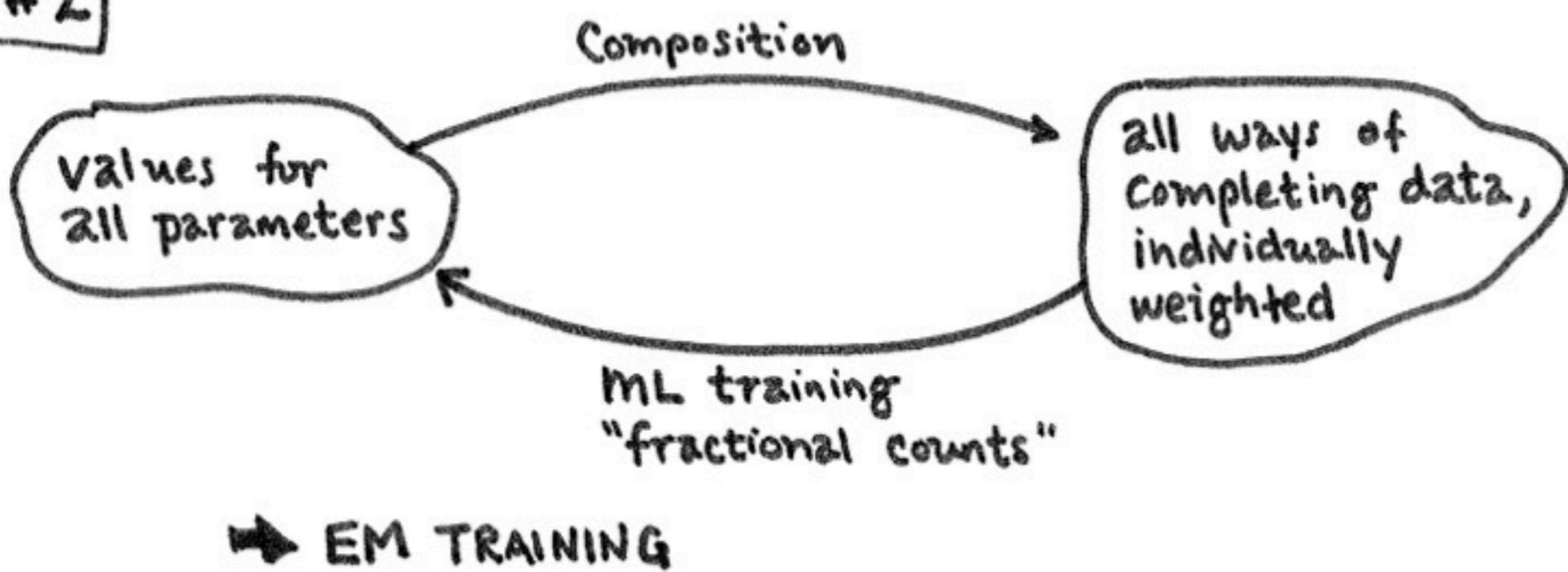
revised  
viterbi:      a a a b a a b a a

# How to Change $m$ ? 2) Soft

Idea #1



Idea #2



# Fractional Counts

- distribution over all possible hallucinated hidden variables

W A I n	W A I n	W A I n	W A I n
W A I N	/ \	\ \	\ \ \
hard-EM counts		0	0
fractional counts	0.333	0.333	0.333
AY  ->	A: 0.333	A I: 0.333	I: 0.333
W  ->	W: 0.667	W A: 0.333	
N  ->	N: 0.667	I N: 0.333	
fractional counts	0.25	0.5	0.25
AY  ->	A I: 0.500	A: 0.250	I: 0.250
W  ->	W: 0.750	W A: 0.250	
N  ->	N: 0.750	I N: 0.250	
eventually	... 0	... I	... 0



# Fractional Counts

- how about

W EH T

W E T O

B IY  
| |\ \  
B I I

B IY  
| \ \  
B I I

- so EM can possibly: (1) learn something correct  
(2) learn something wrong (3) doesn't learn anything
- but with lots of data => likely to learn something good

# EM: slow version (non-DP)

- initialize the conditional prob. table to uniform

- repeat until converged:

- E-step:

$$\begin{array}{ccccccc} & W & A & I & N & & \\ & | & | & / \backslash & & W & A & I & N \\ & & & & & | & | & \backslash & \backslash \\ & & & & & W & A & I & N \\ & & & & & & & & \\ & & & & & z & & z' & z'' \\ & & & & & (z_1 \ z_2 \ z_3) & & & \end{array}$$

- for each training example  $x$  (here:  $(e\dots e, j\dots j)$  pair):
  - for each hidden  $z$ : compute  $p(x, z)$  from the current model
  - $p(x) = \sum_z p(x, z)$ ; [debug: corpus prob  $p(\text{data}) *= p(x)$ ]
  - for each hidden  $z = (z_1 \ z_2 \dots z_n)$ : for each  $i$ :
    - $\text{fraccount}(z_i) += p(x, z)$   $/ p(x)$
- M-step: count-n-divide on fraccounts => new model

# EM: fast version (DP)

- initialize the conditional prob. table to uniform

- repeat until converged:

- E-step:

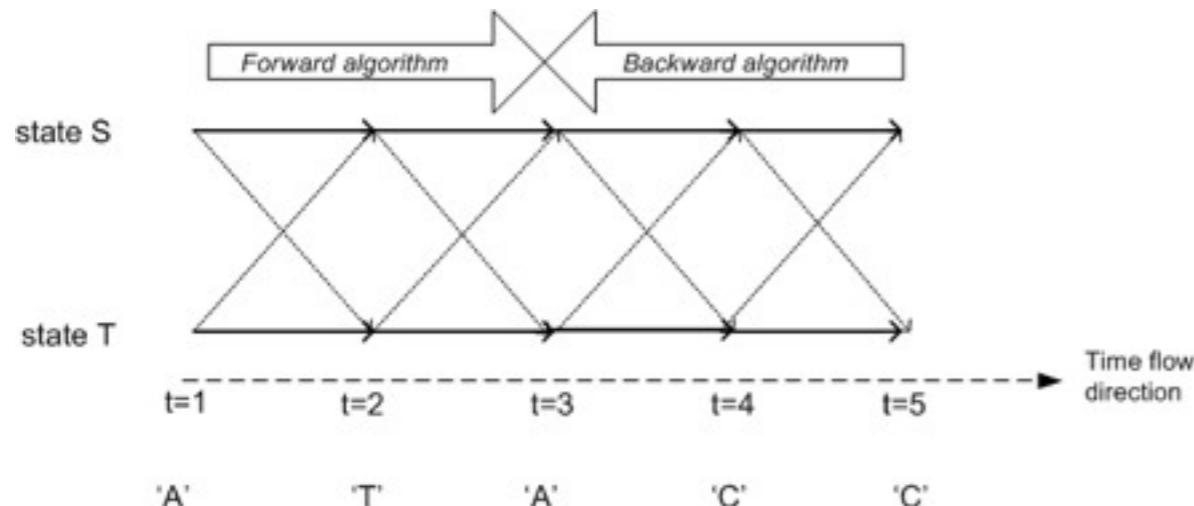
- for each training example  $x$  (here: (e...e, j...j) pair):
  - forward from  $s$  to  $t$ ; note:  $forw[t] = p(x) = \sum_z p(x, z)$
  - backward from  $t$  to  $s$ ; note:  $back[t] = 1$ ;  $back[s] = forw[t]$
  - for each edge  $(u, v)$  in the DP graph with  $label(u, v) = z_i$ 
    - $\text{fraccount}(z_i) += forw[u] * back[v] * prob(u, v) / p(x)$

- M-step: count-n-divide on fraccounts => new model

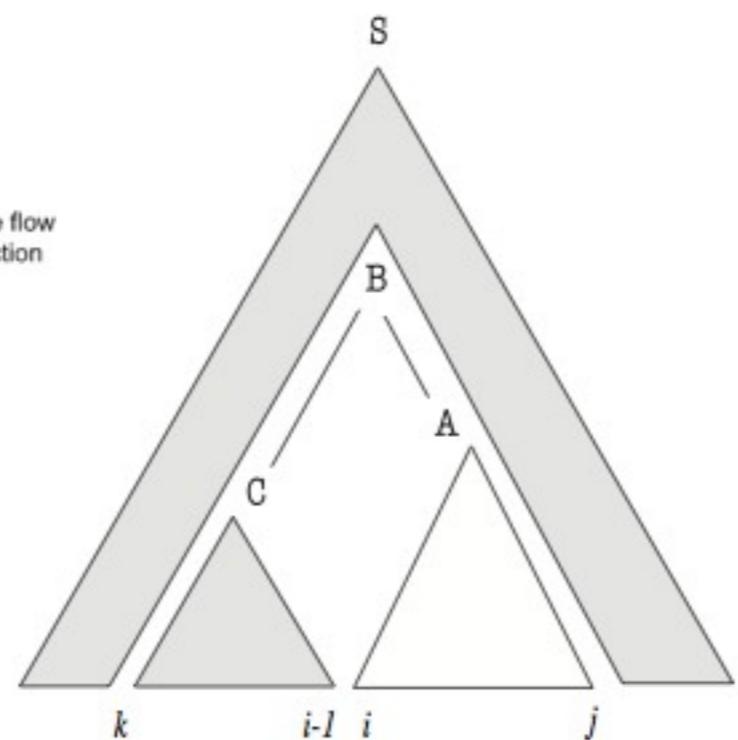
$\sum_{z: (u, v) \in z} p(x, z)$

# How to avoid enumeration?

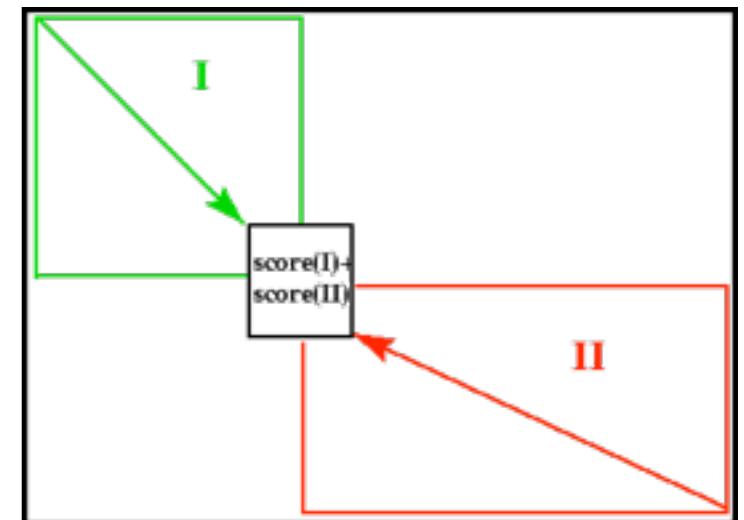
- dynamic programming: the forward-backward algorithm
- forward is just like Viterbi, replacing max by sum
- backward is like reverse Viterbi (also with sum)



POS tagging,  
crypto, ...



inside-outside:  
PCFG, SCFG, ...



alignment,  
edit-distance, ...

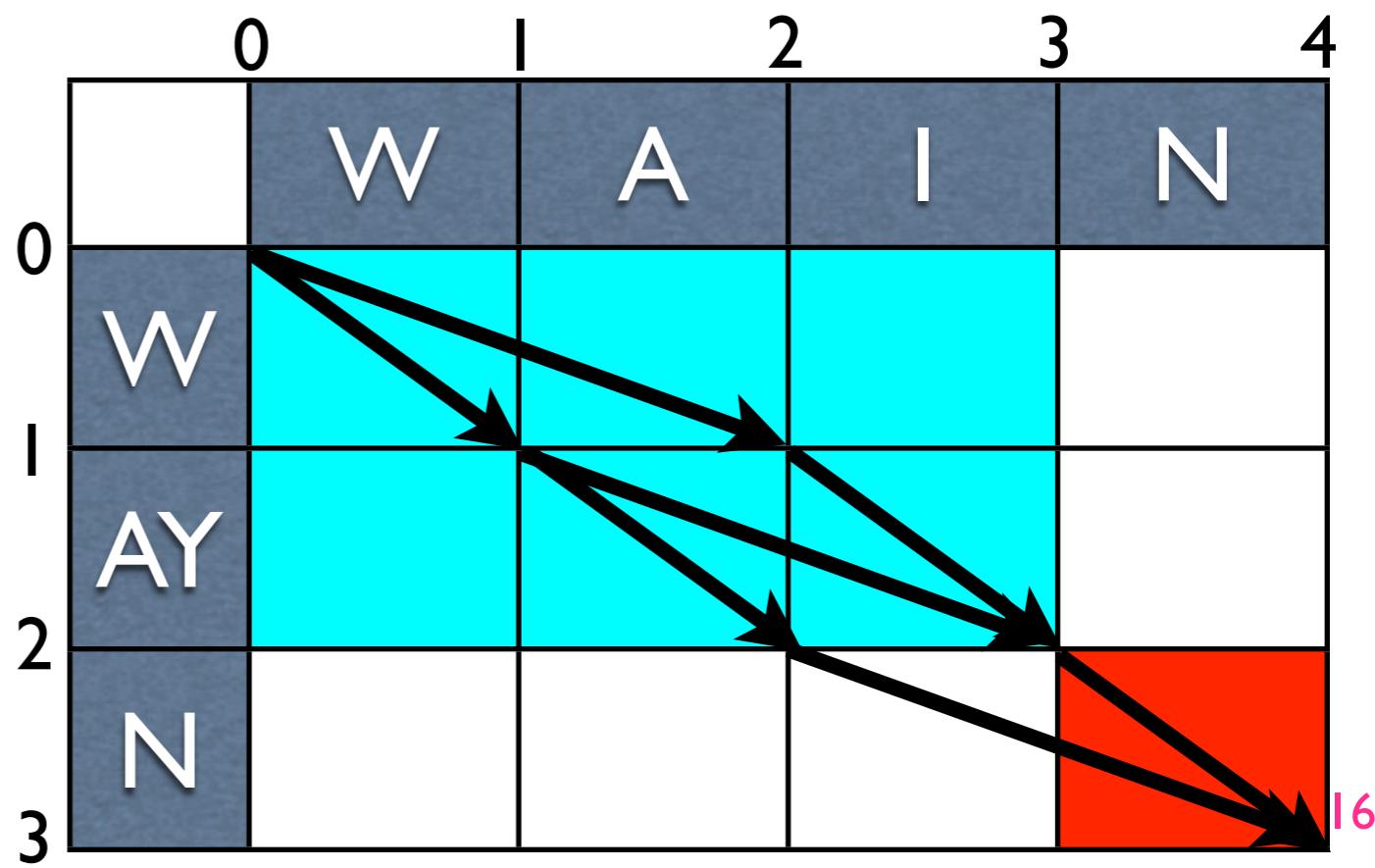
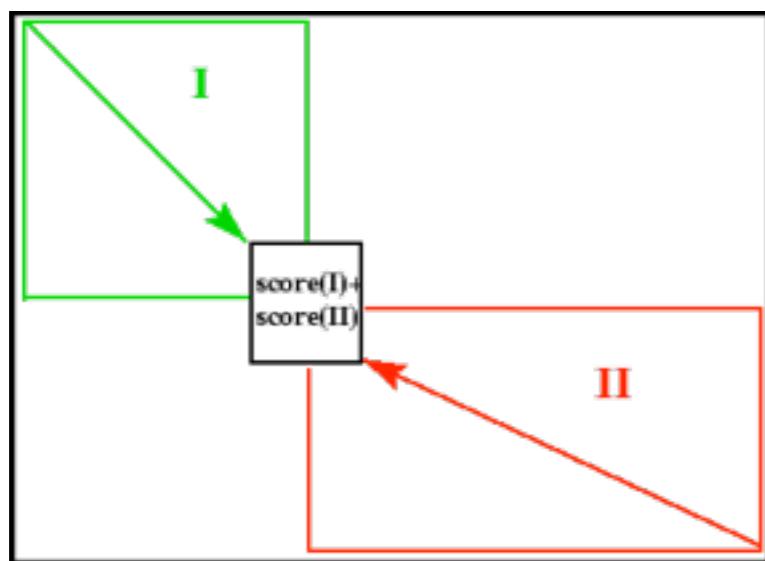
# Example Forward Code

- for HW5. this example shows forward only.

```
n, m = len(eprons), len(jprons)  
forward[0][0] = 1
```

```
for i in xrange(0, n):  
    epron = eprons[i]  
    for j in forward[i]:  
        for k in range(1, min(m-j, 3)+1):  
            jseg = tuple(jprons[j:j+k])  
            score = forward[i][j] * table[epron][jseg]  
            forward[i+1][j+k] += score
```

```
totalprob *= forward[n][m]
```



# Example Forward Code

- for HW5. this example shows forward only.

```
n, m = len(eprons), len(jprons)
```

```
forward[0][0] = 1
```

```
for i in xrange(0, n):
```

```
    epron = eprons[i]
```

```
    for j in forward[i]:
```

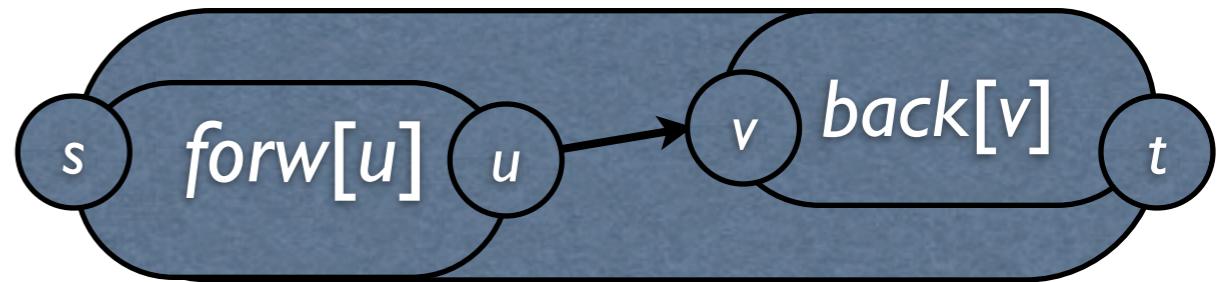
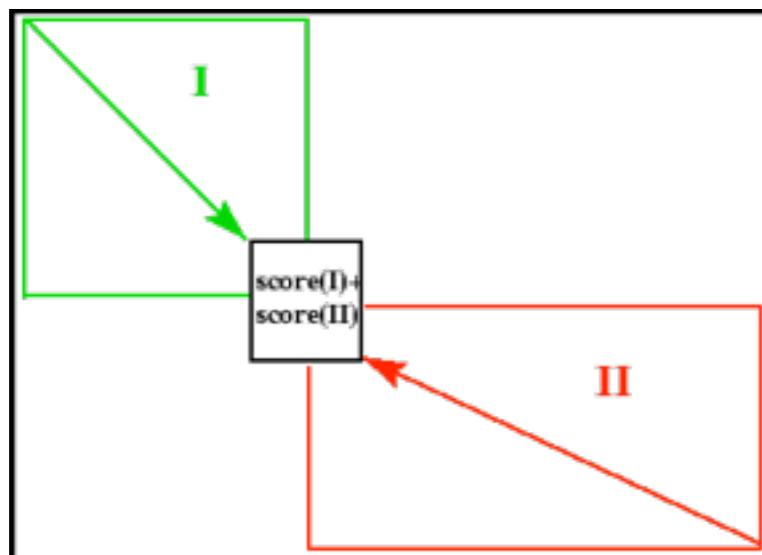
```
        for k in range(1, min(m-j, 3)+1):
```

```
            jseg = tuple(jprons[j:j+k])
```

```
            score = forward[i][j] * table[epron][jseg]
```

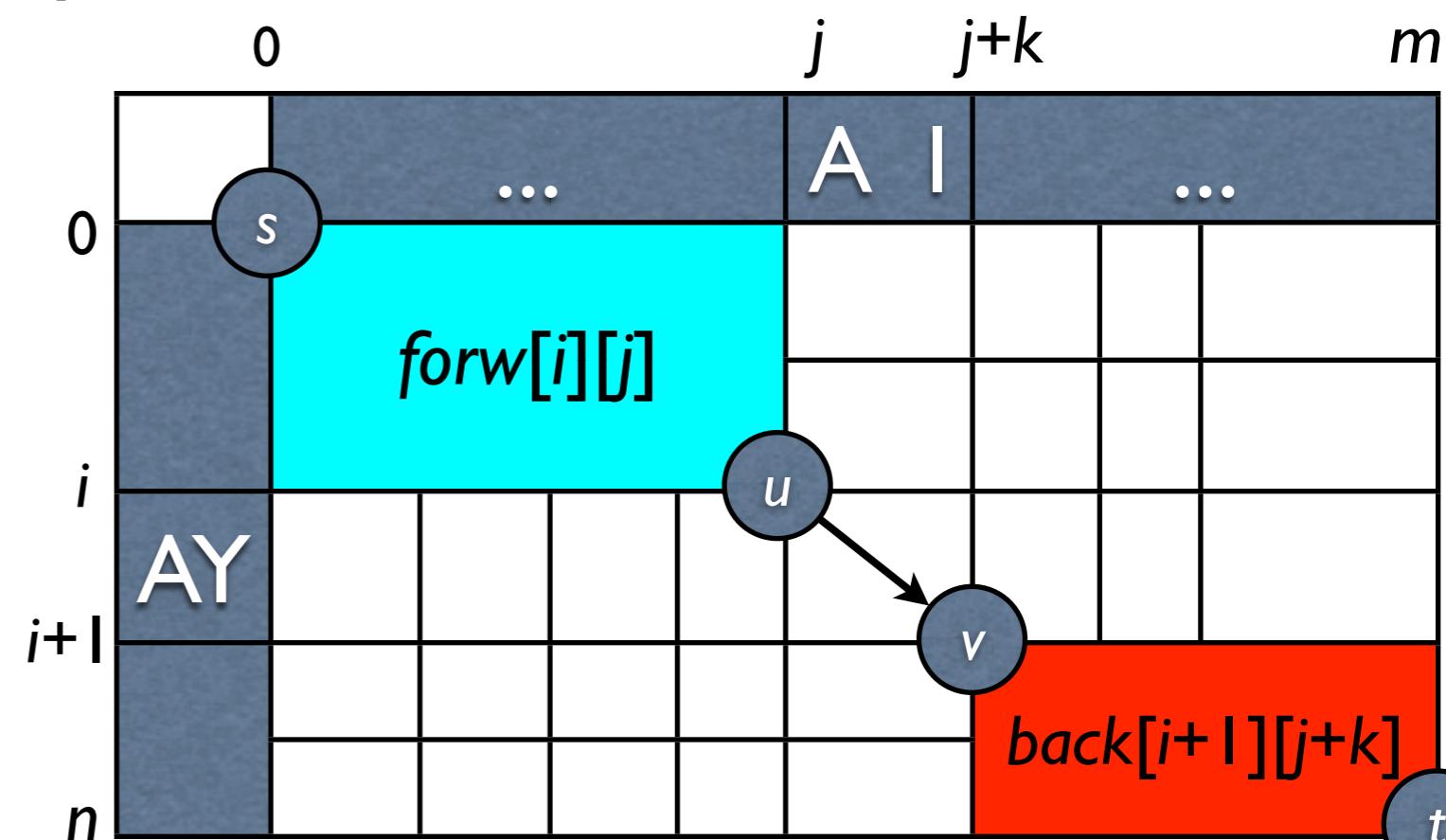
```
            forward[i+1][j+k] += score
```

```
totalprob *= forward[n][m]
```



$forw[s] = back[t] = 1.0$

$forw[t] = back[s] = p(x)$



# EM: fast version (DP)

- initialize the conditional prob. table to uniform

- repeat until converged:

- E-step:

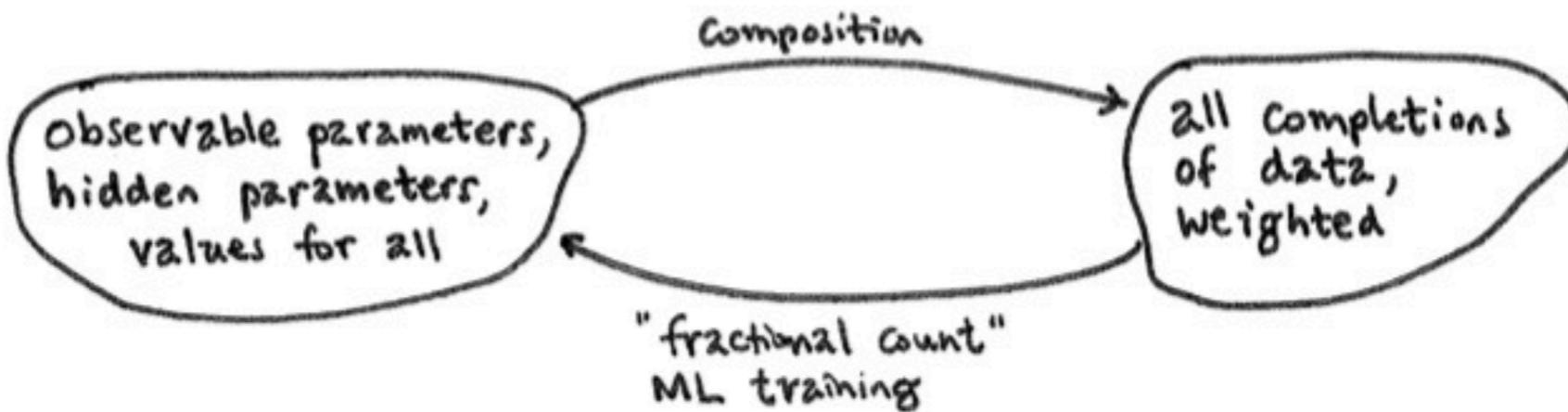
- for each training example  $x$  (here: (e...e, j...j) pair):
  - forward from  $s$  to  $t$ ; note:  $forw[t] = p(x) = \sum_z p(x, z)$
  - backward from  $t$  to  $s$ ; note:  $back[t] = 1$ ;  $back[s] = forw[t]$
  - for each edge  $(u, v)$  in the DP graph with  $label(u, v) = z_i$ 
    - $\text{fraccount}(z_i) += forw[u] * back[v] * prob(u, v) / p(x)$

- M-step: count-n-divide on fraccounts => new model

$\sum_{z: (u, v) \in z} p(x, z)$

# EM

## EM



## example: cryptanalysis

$x_1 \dots x_n$  observed ciphertext

$z_1 \dots z_n$  hidden plaintext

$b(z_j | z_{\text{K}})$  source bigram probabilities OBSERVABLE, FIXED

$t(x_j | z_{\text{K}})$  channel substitution ("encoding") probs HIDDEN

$$P(x_1 \dots x_n, z_1 \dots z_n) = \prod_{i=1}^n b(z_i | z_{i-1}) \cdot t(x_i | z_i) \quad \text{GENERATIVE STORY}$$

$$P(x_1 \dots x_n) = \sum_{z_1 \dots z_n} \prod_{i=1}^n b(z_i | z_{i-1}) \cdot t(x_i | z_i) \quad \text{FORWARD PROCEDURE}$$

$$P(z_1 \dots z_n | x_1 \dots x_n) = \frac{P(x_1 \dots x_n, z_1 \dots z_n)}{P(x_1 \dots x_n)} \quad \text{COND. PROB.}$$

# Why EM increases $p(\text{data})$ iteratively?

$$D = \log p(x; \theta) = \log \sum_z p(x, z; \theta)$$

Note that  $\sum_z p(z|x; \theta_t) = 1$  and  $p(z|x; \theta_t) \geq 0$  for all  $z$ . Therefore  $D$  is the logarithm of a weighted sum, so we can apply Jensen's inequality, which says  $\log \sum_j w_j v_j \geq \sum_j w_j \log v_j$ , given  $\sum_j w_j = 1$  and each  $w_j \geq 0$ . Here, we let the sum range over the values  $z$  of  $Z$ , with the weight  $w_j$  being  $p(z|x; \theta_t)$ . We get

$$D \geq E = \sum_z p(z|x; \theta_t) \log \frac{p(x, z; \theta)}{p(z|x; \theta_t)}.$$

Separating the fraction inside the logarithm to obtain two sums gives

$$E = \left( \sum_z p(z|x; \theta_t) \log p(x, z; \theta) \right) - \left( \sum_z p(z|x; \theta_t) \log p(z|x; \theta_t) \right).$$

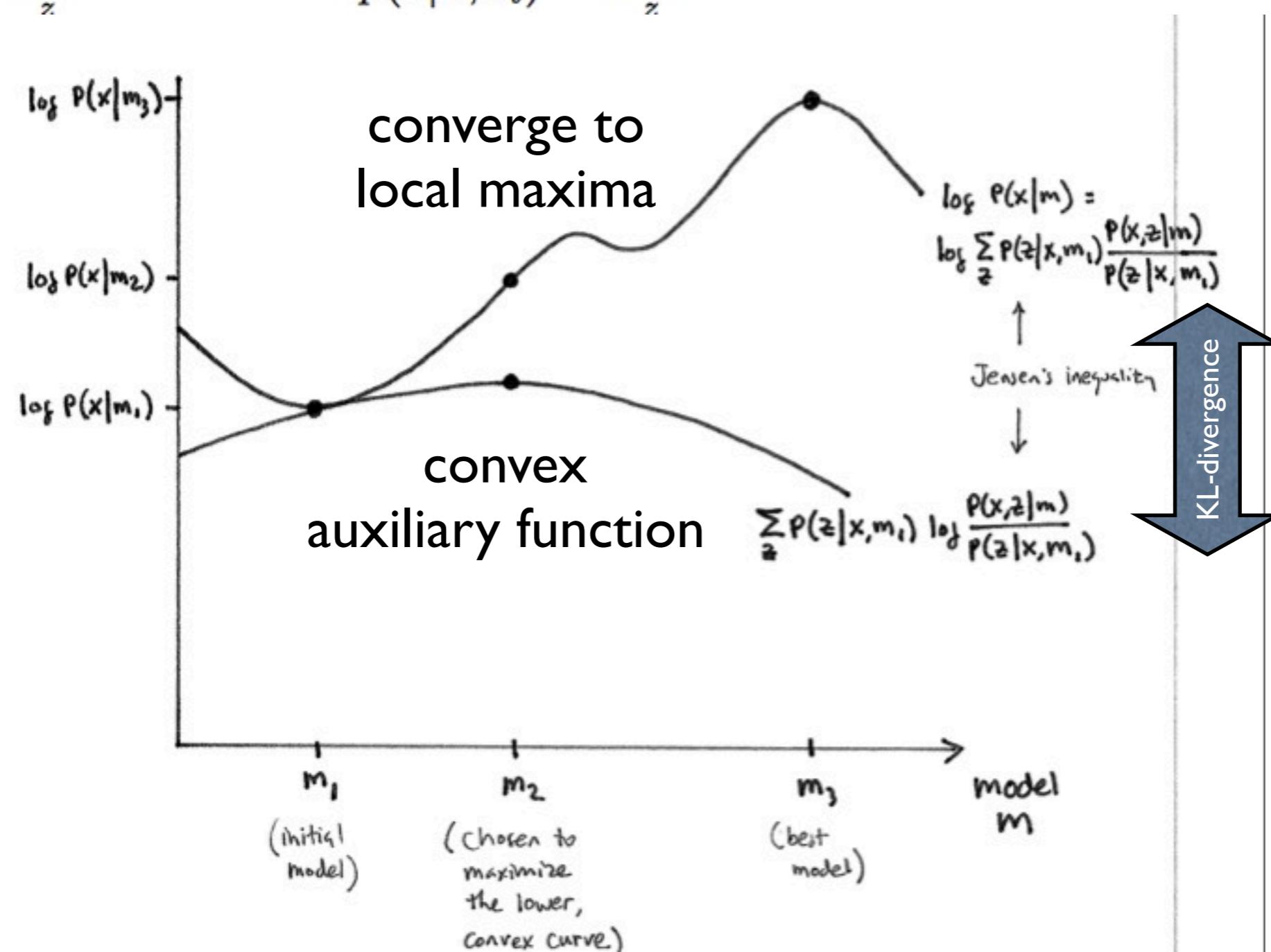
Since  $E \leq D$  and we want to maximize  $D$ , consider maximizing  $E$ . The weights  $p(z|x; \theta_t)$  do not depend on  $\theta$ , so we only need to maximize the first sum, which is

$$\sum_z p(z|x; \theta_t) \log p(x, z; \theta).$$

# Why EM increases $p(\text{data})$ iteratively?

How do we know that maximizing  $E$  actually leads to an improvement in the likelihood? With  $\theta = \theta_t$ ,

$$E = \sum_z p(z|x; \theta_t) \log \frac{p(x, z; \theta_t)}{p(z|x; \theta_t)} = \sum_z p(z|x; \theta_t) \log p(x; \theta_t) = \log p(x; \theta_t)$$



# How to maximize the auxiliary?

$$\sum_z p(z|x; \theta_t) \log p(x, z; \theta).$$

In general, the E-step of an EM algorithm is to compute  $p(z|x; \theta_t)$  for all  $z$ . The M-step is then to find  $\theta$  to maximize  $\sum_z p(z|x; \theta_t) \log p(x, z; \theta)$ .

W A I N	W A I N	W A I N
/ \	\ \ \	\ \ \ \ \
W A I N	W A I N	W A I N
$p(z x) = 0.5$	$p(z' x) = 0.3$	$p(z'' x) = 0.2$

just count-n-divide on  
the fractional data!  
(as if MLE on complete data)

W A I N	W A I N	W A I N
/ \	\ \ \	\ \ \ \ \
W A I N	W A I N	W A I N
$5x$	$3x$	$2x$

