

Lecture 3: Information Theory Continues



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Review:

Shannon's Information Theory

The
Claude Shannon: ~~A~~ Mathematical Theory of Communication

Bell System Technical Journal, 1948

- Shannon's measure of information is the number of bits to represent the amount of uncertainty (randomness) in a data source, and is defined as **entropy**

$$H = -\sum_{i=1}^n p_i \log(p_i)$$

Where there are n symbols $1, 2, \dots, n$, each with probability of occurrence of p_i

Entropy: Three properties

1. It can be shown that $0 < H < \log N$.
2. **Maximum entropy** ($H = \log N$) is reached when all symbols are **equiprobable**, i.e.,
 $p_i = 1/N$.
3. The difference $\log N - H$ is called the **redundancy** of the source.

Joint Information

- X and Y are random variables.
- X and Y can have n and m possibilities, respectively. Then, the joint information is defined as:

$$H(X, Y) = - \sum_{i=1}^n \sum_{j=1}^m r(x_i, y_j) \log(r(x_i, y_j))$$

- $r(x, y)$ is the joint probability of x and y.
- Why this definition?

Conditional Information

- X and Y are random variables.
- X and Y can have n and m possibilities, respectively. Then, the conditional information is defined as:

$$H(Y | X) = - \sum_{i=1}^n \sum_{j=1}^m r(x_i, y_j) \log(q(y_j | x_i))$$

- $q(y|x)$ is the conditional probability.
- Why this definition?

Conditional Information

Properties of Conditional Information:

1. $H(Y | X) \geq 0$
2. $H(Y | X) \leq H(Y)$ with equality if X and Y are independent.
3. $H(X, Y) = H(X) + H(Y | X) = H(Y) + H(X | Y)$

Mutual Information

- X and Y are random variables.
- X and Y can have n and m possibilities, respectively.
Then, the mutual information is defined as:

$$I(X, Y) = H(Y) - H(Y | X) = \sum_{i=1}^n \sum_{j=1}^m r(x_i, y_j) \log \left[\frac{r(x_i, y_j)}{q(y_j) p(x_i)} \right]$$

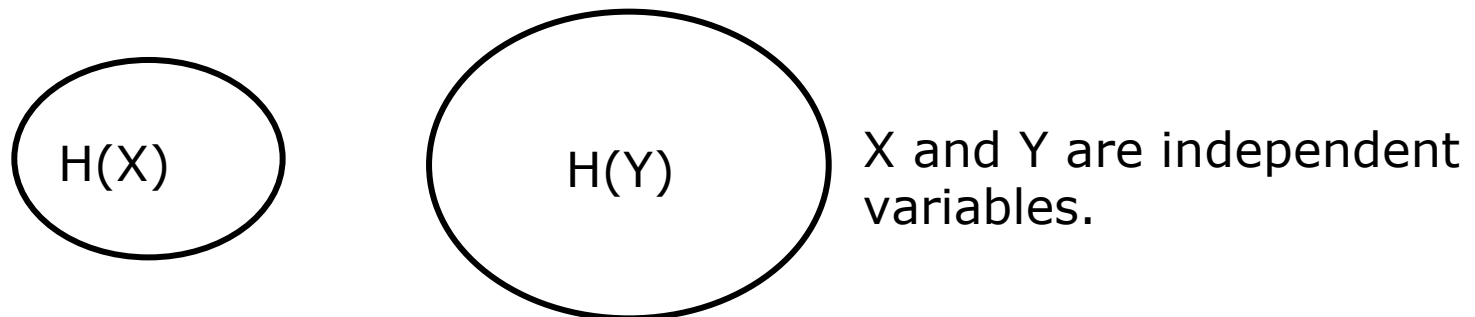
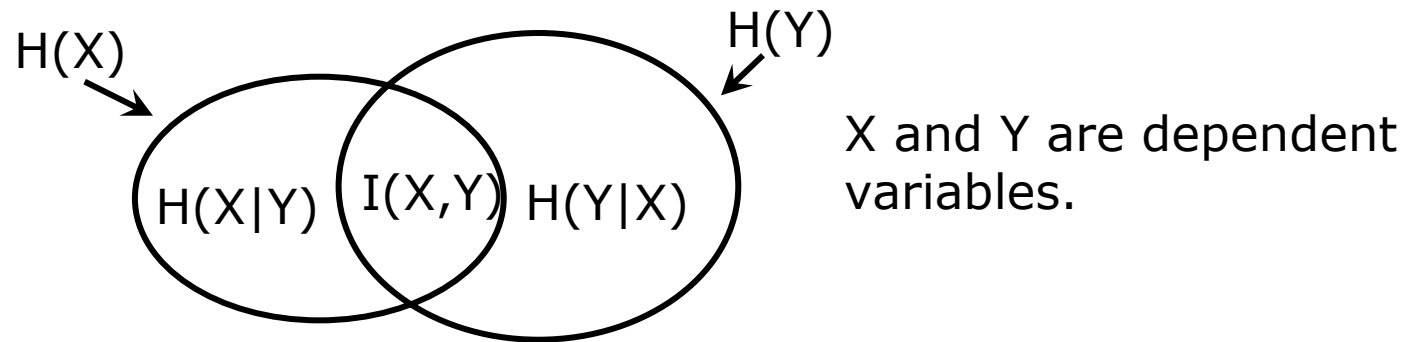
- Why this definition?

Mutual Information

- Properties of Mutual Information:

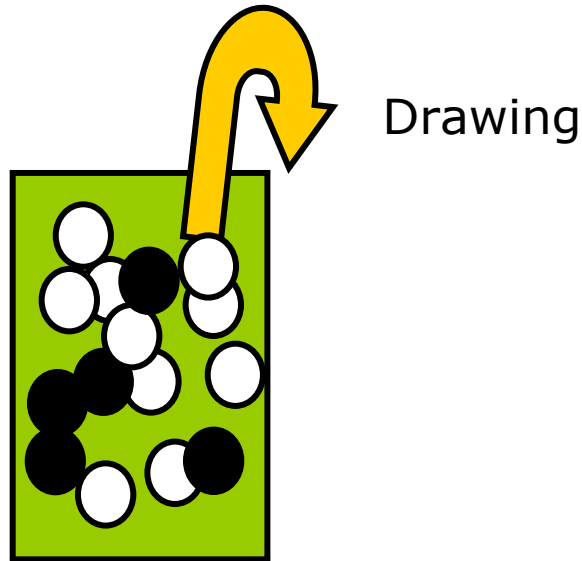
$$I(X, Y) = I(Y, X)$$

Relationship among entropy, conditional, and mutual information



Example:

- A vase contains **5** black balls and **10** white balls. Experiment x involves the random drawing of a ball, without being replaced in the vase. Experiment Y involves random drawing of the second ball.
- **5** black balls
- **10** white balls

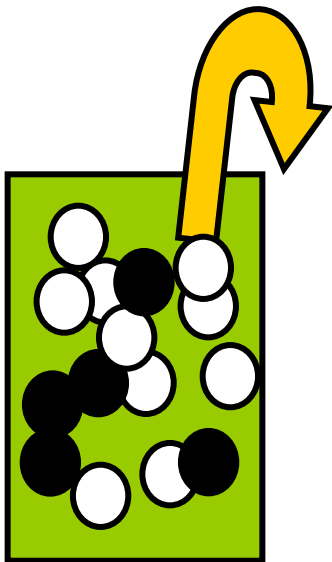


Example: Entropy

- How much uncertainty (information) does experiment X contain?

$$P(\text{black}_X) = 1/3, P(\text{white}_X) = 2/3$$

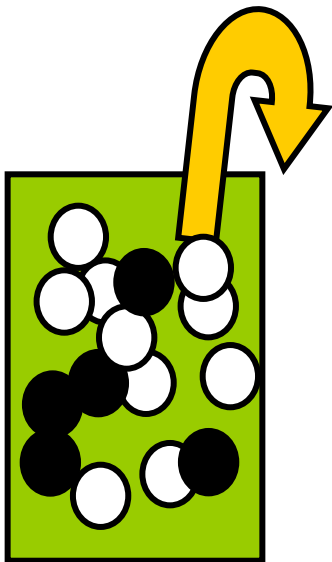
$$H(X) = -(1/3)\log(1/3) - (2/3)\log(2/3) = 0.92 \text{ bit}$$



Drawing

Example:

- How much uncertainty (information) in experiment Y given that the ball in experiment X is white?



Drawing

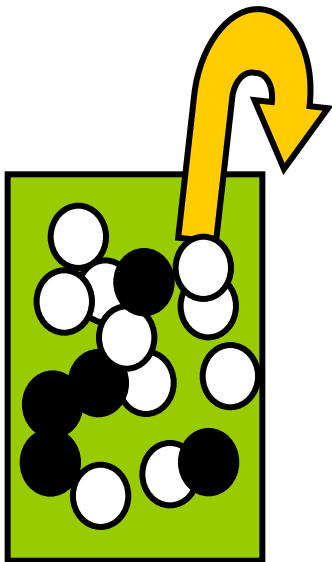
$$P(\text{black}_Y|\text{white}_X) = 5/14$$

$$P(\text{white}_Y|\text{white}_X) = 9/14$$

$$H(Y|\text{white}_X) = -(5/14)\log(5/14) - (9/14)\log(9/14) = .94 \text{ bit}$$

Example:

- How much uncertainty (information) in experiment Y given that the ball in experiment X is black?



Drawing

$$P(\text{black}_Y|\text{black}_X) = 4/14 = 2/7$$

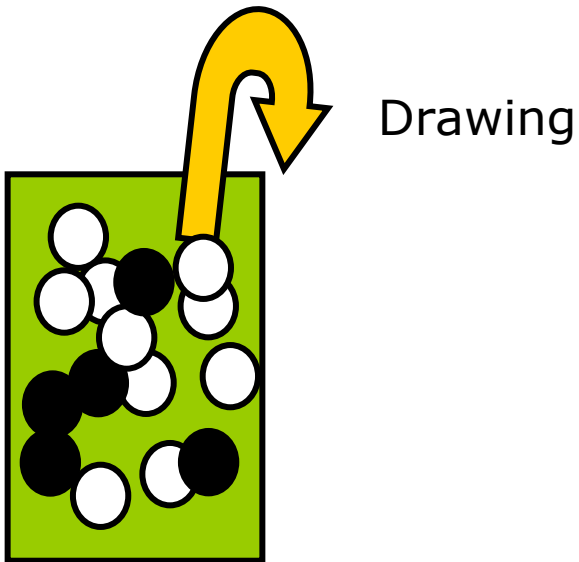
$$P(\text{white}_Y|\text{black}_X) = 10/14 = 5/7$$

$$H(Y|\text{black}_X) = -(2/7)\log(2/7) - (5/7)\log(5/7) = 0.86 \text{ bit}$$

Example:

- How much uncertainty does experiment Y contain?

$$\begin{aligned} H(Y) &= P(\text{black}_X) * H(Y|\text{black}_X) + P(\text{white}_X) * H(Y|\text{white}_X) \\ &= (1/3)(0.86) + (2/3)(0.94) = 0.91 \text{ bit} \end{aligned}$$



Formal Derivation of Entropy

- Why do we have

$$H = - \sum_{i=1}^n p_i \log(p_i) ?$$

Axiomatic Foundations

- Assuming that information measure should satisfy the three following requirements (Chaundy and McLeod (1960)):
 1. If all outcomes are split up into groups, then all the values of H for the various groups, multiplied by the statistical weights, should lead to the overall H .
 2. H should be continuous in p_i .
 3. If all p_i 's are equal, i.e. for all i , ($p_i = 1/n$), then H will increase monotonically as a function of n . That means the uncertainty will increase for an increasing number of equal probabilities.

Derivation of Entropy

- Theorem: The only function that satisfy the three requirements above is

$$H = -K \sum_{i=1}^n p_i \log(p_i)$$

- Proof:

Summary

- History of information theory.
- Information theoretical entities
 - Information, self-information, entropy, conditional information, joint information, mutual information.
- Derivation of $H = - \sum p_i \log p_i$